# The Ordinal Numbers and Transfinite Induction

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## 1 Intuition behind Cardinals and Ordinals

2 Construction of Ordinal Numbers

- 3 Ordinal Number Arithmetic
- Transfinite Recursion and Induction

## 1 Intuition behind Cardinals and Ordinals

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Transfinite Recursion and Induction

A cardinal number represents the *size* of a set.

The cardinality of a set S is denoted by |S|.

# Examples $$\begin{split} |\varnothing| &= 0 \\ |\{0,1,2\}| &= 3 \\ |\mathbb{N}| &= |\mathbb{Q}| = \aleph_0, \text{ the first transfinite cardinal} \\ |\mathbb{R}| &= \mathfrak{c}, \text{ the cardinality of the continuum, strictly larger than } \aleph_0 \end{split}$$

Intuitively, one would think to define the cardinal numbers as equivalence classes of sets, where  $S \sim T$  if and only if there exists a bijection  $f: S \rightarrow T$ . This is how Bertrand Russell defined a "number."

#### Definition

A set S is said to be **well-ordered** if S is totally ordered and every nonempty subset of S has a least member.

The canonical example of a well-ordered set is  $\mathbb{N}$ .

If we assume the Axiom of Choice (which we will in this talk), we can prove the Well-Ordering Theorem.

#### Well-Ordering Theorem

Every set can be well-ordered.

#### Definition

Two well-ordered sets  $(S, \leq)$  and  $(T, \leq')$  are said to be **order isomorphic** if there exists a bijection  $f : S \to T$  such that  $s \leq s'$  if and only if  $f(s) \leq' f(s')$ . In this case, f is said to be an **order isomorphism**.

#### Examples

• If  $S = \{1,2\}$  and  $T = \{3,15\}$  under the standard ordering, then

 $f: S \rightarrow T$  given by f(1) = 3 and f(2) = 15 is an order isomorphism.

•  $g:\mathbb{N}\to\mathbb{N}_+$  by g(n)=n+1 is an order isomorphism under the standard ordering.

• Let  $\mathbb{N} \cup \{\omega\}$  be given the ordering that  $n < \omega$  for all natural numbers n. Then there is no order isomorphism between  $\mathbb{N}$  and  $\mathbb{N} \cup \{\omega\}$ .

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An ordinal number can be thought of as the *position* of an element in a well-ordered set.

#### Example

Let  $\mathbb{N} \cup \{\omega\}$  have the same ordering as before. The 1st element is 0. The 14th element is 13. The *n*th element is n-1. The  $(\omega + 1)$ st element is  $\omega$ .

An ordinal number can also be thought of as the *order type* of a well-ordered set (the position of the "last" element, if one exists).

#### Examples

The order type of  $\varnothing$  is 0. The order type of  $\{0, 1, 2\}$  is 3. The order type of  $\mathbb{N}$  is  $\omega$ .

# Problems with Equivalence Classes

Like before, we could think of ordinal numbers as equivalence classes of well-ordered sets, where  $(S, \leq) \sim (T, \leq')$  if and only if  $(S, \leq)$  is order isomorphic to  $(T, \leq')$ .

This view of the ordinal numbers and the corresponding view of the cardinal numbers does *not* work in ZFC.

It turns out that bijection classes of sets are too big to be sets. Similarly, order isomorphism classes of well-ordered sets are too big to be sets. This is related to Russell's Paradox. A collection C is too big to be a set if

 ${\mathcal C}$  a set  $\Rightarrow$  the set of all sets exists

There are a few ways to prove that a collection C is too big to be a set.

$$\mathcal{C} \text{ a set} \Rightarrow \mathcal{C} \in \mathcal{C}$$

 $S \hookrightarrow \mathcal{C}$  for all sets S

 $\mathcal D$  is too big to be a set and  $\mathcal D \hookrightarrow \mathcal C$ 

## Intuition behind Cardinals and Ordinals

## 2 Construction of Ordinal Numbers

- 3 Ordinal Number Arithmetic
- 4 Transfinite Recursion and Induction

The von Neumann construction of the natural numbers is as follows:  $\mathbf{0} := \varnothing$ 

$$n+1 := n \cup \{n\} = \{0, 1, 2, ..., n\}$$

We can extend this construction to make a definition for every ordinal number.

An ordinal number  $\alpha$  is defined to be the set of all ordinal numbers less than  $\alpha.$ 

| Examples  |  |
|---|--|
| 0 := arnothing  |  |
| $5 \coloneqq \{0, 1, 2, 3, 4\}$                               |  |
| $\omega\coloneqq \mathbb{N}=\{0,1,2,\}$                       |  |
| $\omega+1\coloneqq\mathbb{N}\cup\{\omega\}=\{0,1,2,;\omega\}$ |  |

#### Definition

An ordinal number  $\beta$  is a **successor ordinal** if there exists an ordinal number  $\alpha$  such that  $\beta = \alpha + 1 = \{0, 1, ..., \alpha\}$ . An ordinal number  $\lambda$  is a **limit ordinal** if  $\lambda > 0$  and for all ordinals  $\alpha$ ,  $\lambda \neq \alpha + 1$ .

Successor ordinals are named as such because they are the immediate successor of some other ordinal. Limit ordinals are named as such because they are the limit of the sequence of ordinals less then them without being a successor ordinal.

#### Examples

All natural numbers n > 0 are successor ordinals.  $\omega + n$  is a successor ordinal for all natural numbers n > 0.  $\omega$  is the smallest limit ordinal.

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Every well-ordered set is order isomorphic to exactly one ordinal number (and the isomorphism is unique!). As such, we make the following definition:

#### Definition

The order type of a well-ordered set  $(S, \leq)$  is the unique ordinal number which is order isomorphic to  $(S, \leq)$ . Denote the order type of  $(S, \leq)$  as **Ord**  $(S, \leq)$ .

These facts allow us to see that this new definition of ordinal number fits in line with our intuitive notion of an ordinal number from before.

# **Cardinal Numbers**

Even though every set can be well-ordered (Well-Ordering Theorem), for infinite sets, the order type can depend on the specific ordering chosen.

#### Examples

**Ord**  $\mathbb{N} = \omega$  under the standard ordering.

If we choose an ordering  $\leq'$  on  $\mathbb{N}$  so that 0 is greater than every other natural number, but the other natural numbers have the standard ordering, then  $\mathbf{Ord} (\mathbb{N}, \leq') = \omega + 1 \neq \omega$ .

With this in mind, we can define cardinality in the following way.

#### Definition

Let S be a set. The **cardinality** of S (denoted |S|) is the minimum order type of S under all possible well-orderings. i.e.,  $|S| = \min_{\leq} \operatorname{Ord}(S, \leq)$ .

We skip the proofs for time purposes, but one can prove that there exists a bijection between S and T iff |S| = |T|.

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## Intuition behind Cardinals and Ordinals

2 Construction of Ordinal Numbers

### Ordinal Number Arithmetic

Transfinite Recursion and Induction

#### Definition

Let  $\alpha$  and  $\beta$  be ordinal numbers. Consider the set  $\alpha \amalg \beta$ . Define an well-order  $\leq$  on  $\alpha \amalg \beta$  by preserving the original orders on  $\alpha$  and  $\beta$  themselves and having the additional criterion that all elements of  $\beta$  are greater than all elements of  $\alpha$ . Then  $\alpha + \beta := \operatorname{Ord} (\alpha \amalg \beta, \leq)$ .

#### Examples

$$\begin{array}{l} 2+3 = \mathbf{Ord} \left( \{0,1\} \amalg \{0,1,2\} \right) = \mathbf{Ord} \left\{ 0,1,0',1',2' \right\} = 5 \\ \omega + 1 = \mathbf{Ord} \left( \mathbb{N} \amalg \{0\} \right) = \mathbf{Ord} \left\{ 0,1,2,...;0' \right\} = \omega + 1 \\ 1+\omega = \mathbf{Ord} \left( \{0\} \amalg \mathbb{N} \right) = \mathbf{Ord} \left\{ 0,0',1',2',... \right\} = \omega \end{array}$$

The example of  $\omega+1$  justifies our choice of notation for successors. The example of  $1+\omega$  shows that addition of ordinal numbers is not commutative!

#### Definition

Let  $\alpha$  and  $\beta$  be ordinal numbers. Consider the set  $\alpha \times \beta$ . Define an well-order < on  $\alpha \times \beta$  by  $(\alpha', \beta') < (\alpha'', \beta'')$  if and only if  $\beta' < \beta''$  or  $\beta' = \beta''$  and  $\alpha' < \alpha''$ . For convenience, we will denote  $(\alpha', \beta')$  by  $\alpha'_{\beta'}$ . Then  $\alpha\beta := \operatorname{Ord} (\alpha \times \beta, <)$ .

#### Examples

$$\begin{array}{l} 2 \cdot 3 = \mathbf{Ord} \left( \{0, 1\} \times \{0, 1, 2\} \right) = \mathbf{Ord} \left\{ 0_0, 1_0, 0_1, 1_1, 0_2, 1_2 \right\} = 6 \\ \omega 2 = \mathbf{Ord} \left( \mathbb{N} \times \{0, 1\} \right) = \mathbf{Ord} \left\{ 0_0, 1_0, 2_0, ...; 0_1, 1_1, 2_1, ... \right\} = \omega + \omega \\ 2\omega = \mathbf{Ord} \left( \{0, 1\} \times \mathbb{N} \right) = \mathbf{Ord} \left\{ 0_0, 1_0, 0_1, 1_1, 0_2, 1_2 ... \right\} = \omega \end{array}$$

Again, we see that multiplication is not commutative!

So far, we have seen some operations and shown that they are quite complicated. For time purposes, we will skip ordinal exponentiation. In ordinal exponentiation, we get even weirder results. For example, there are types of ordinal numbers called " $\varepsilon$ -numbers" which satisfy the equation

$$\varepsilon=\omega^\varepsilon$$

The first  $\varepsilon$ -number is  $\varepsilon_0 = \omega^{\omega^{\omega^*}}$ 

In this, we are dealing with unfathomably large numbers, but we've barely scratched the surface.

The collection of ordinal numbers is too big to be a set. Indeed, if it were a set, it would itself be an ordinal number, so it would have to contain itself as an element.

Overall, the ordinal numbers are very complicated, but they can be quite helpful when it comes to transfinite recursion and induction.

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2 Construction of Ordinal Numbers

3 Ordinal Number Arithmetic

#### Transfinite Recursion and Induction

Recursion and Induction are similar processes.

Recursion *defines* a sequence of objects.

Induction *proves* a property is true for a sequence of objects.

Normally, we work with recursion and induction for only the natural numbers, giving us sequences of order type  $\omega$ .

When we have sequences of order type larger than  $\omega$ , we use *transfinite* recursion and induction, since our processes use transfinite ordinal numbers.

To define things recursively, we use the following process:

#### Transfinite Recursion

1. Base Case: Define an object  $S_0$ .

2. Successor Step: For all ordinal numbers  $\alpha$ , define  $S_{\alpha+1}$  using the definition of  $S_{\alpha}$ .

3. Limit Step: For all limit ordinals  $\lambda$ , define  $S_{\lambda}$  using the definition of  $S_{\alpha}$  for all ordinals  $\alpha < \lambda$ .

Consider  $\mathbb{R}$  as a  $\mathbb{Q}$ -vector space. By Zorn's Lemma, we know this vector space must have a basis. Here, we will use transfinite recursion and induction to prove that this vector space has a basis.

Base Case: Let  $B_0 = \emptyset$ .

Successor Case: For all ordinals  $\alpha$ , choose  $v_{\alpha+1}$  to be an element of  $\mathbb{R}$  – Span  $B_{\alpha}$ . Then, define  $B_{\alpha+1} = B_{\alpha} \cup \{v_{\alpha+1}\}$ .

Limit Case: For all limit ordinals  $\lambda$ , define  $B_{\lambda} = \bigcup_{\alpha < \lambda} B_{\alpha}$ .

Since  $|\mathbb{R}| < |\wp(\mathbb{R})| = \gamma$  for some ordinal  $\gamma$ , it follows that this process will terminate at some ordinal  $\beta$ , where  $\mathbb{R} - \text{Span } B_{\beta} = \emptyset$ .

Claim:  $B_{\beta}$  is a basis for  $\mathbb{R}$  as a  $\mathbb{Q}$ -vector space.

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To prove things inductively, we use the following process:

#### Transfinite Induction

Let *P* be a property and  $P(\alpha)$  denote that *P* is true for the ordinal  $\alpha$ . The following is the method of transfinite induction to prove that *P* is true for all ordinal numbers:

- 1. Base Case: Prove P(0).
- 2. Successor Step: Assuming  $P(\alpha)$  for arbitrary  $\alpha$ , prove  $P(\alpha + 1)$ .

3. Limit Step: For all limit ordinals  $\lambda$ , assuming  $P(\alpha)$  for all ordinals  $\alpha < \lambda$ , prove  $P(\lambda)$ .

To show that  $B_{\beta}$  is a basis, it suffices to show that  $B_{\beta}$  is linearly independent. Afterall,  $\mathbb{R}$  – Span  $B_{\beta} = \emptyset$ , so  $B_{\beta}$  spans  $\mathbb{R}$ . To do this, it suffices to show that  $B_{\alpha}$  is linearly independent for all ordinals  $\alpha$ .

<u>Base case</u>:  $B_0 = \emptyset$  is vacuously linearly independent.

<u>Successor case</u>: Suppose that for some ordinal  $\alpha$ ,  $B_{\alpha}$  is linearly independent, but  $B_{\alpha+1}$  is not. Now,  $B_{\alpha+1} = B_{\alpha} \cup \{v_{\alpha+1}\}$ , so it follows that  $0 = c_1w_1 + \ldots + c_nw_n + c_{n+1}v_{\alpha+1}$ . We can then write

$$v_{\alpha+1} = -c_{n+1}^{-1}c_1w_1 - \dots - c_{n+1}^{-1}c_nw_n$$

So  $v_{\alpha+1} \in \text{Span } B_{\alpha}$ , which contradicts the fact that  $v_{\alpha+1} \in \mathbb{R} - \text{Span } B_{\alpha}$ . Thus, when  $B_{\alpha}$  is linearly independent,  $B_{\alpha+1}$  is linearly independent.

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# Proof that $B_{\beta}$ is a basis.

<u>Limit case</u>: Suppose that  $\lambda$  is a limit ordinal. Suppose that  $B_{\alpha}$  is linearly independent for all  $\alpha < \lambda$ .

Notice that, by definition,  $B_{\gamma} \subset B_{\delta}$  for all ordinals  $\gamma < \delta$ . Hence,  $\{B_{\alpha}\}_{\alpha < \lambda}$  is actually a chain of sets.

Now, suppose that  $B_{\lambda}$  is linearly dependent. Then there exists a finite set  $\{w_1, ..., w_n\} \subset B_{\lambda}$  which is linearly dependent. By definition,  $B_{\lambda} = \bigcup_{\alpha < \lambda} B_{\alpha}$ , so it follows that there exist ordinals  $\alpha_1, ..., \alpha_n < \lambda$  with  $w_i \in B_{\alpha_i}$ .

Since  $\{B_{\alpha}\}_{\alpha<\lambda}$  is a chain, without loss of generality, it follows that  $B_{\alpha_1} \subset ... \subset B_{\alpha_n}$ . In particular,  $\{w_1, ..., w_n\} \subset B_{\alpha_n}$ , which contradicts the assumption that  $B_{\alpha_n}$  is linearly independent (since  $\alpha_n < \lambda$ ).

Therefore,  $B_{\lambda}$  is linearly independent whenever  $B_{\alpha}$  is linearly independent for all  $\alpha < \lambda$ .  $\Box$ 

# Every well-ordered set has an order type

We will now give another example of transfinite induction to prove that every well-ordered set has an order type. To do this, we will prove that every well-ordered set is uniquely order isomorphic to some ordinal.

Let  $(S, \leq)$  be a well-ordered set. We will begin by using transfinite recursion to construct a sequence of subsets of  $(S, \leq)$ .

<u>Base case</u>: Let  $S_0 = \emptyset$ .

<u>Successor case</u>: For all ordinals  $\alpha$ , define  $x_{\alpha}$  to be the least element of  $S - S_{\alpha}$ , provided that  $S - S_{\alpha} \neq \emptyset$ . Then  $S_{\alpha+1} = S_{\alpha} \cup \{x_{\alpha}\}$ .

<u>Limit case</u>: For all limit ordinals  $\lambda$ , define  $S_{\lambda} = \bigcup_{\alpha < \lambda} S_{\alpha}$ .

This process must terminate for some ordinal  $\beta$  (where  $S - S_{\beta} = \emptyset$ ). If not, then the ordinal numbers can be embedded into S, which makes S too big to be a set, contradicting the fact that S is a set.

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## Every well-ordered set has an order type

We prove that  $S_{\alpha}$  is uniquely order isomorphic to  $\alpha$  for all ordinals  $\alpha \leq \beta$ .

<u>Base case</u>:  $S_0 = \emptyset$  is vacuously order isomorphic to  $0 = \emptyset$  by the empty function, which is unique.

<u>Successor case</u>: Suppose  $S_{\alpha}$  is uniquely order isomorphic to  $\alpha$  for some ordinal  $\alpha < \beta$ . i.e.,  $f_{\alpha} : S_{\alpha} \to \alpha$  is the unique order isomorphism.

Now, by definition,  $S_{\alpha+1} = S_{\alpha} \cup \{x_{\alpha}\}$ , where  $x_{\alpha}$  is the smallest element of  $S - S_{\alpha}$ . It follows that  $x_{\alpha}$  is larger than all elements of  $S_{\alpha}$ . So we establish the function  $f_{\alpha+1} : S_{\alpha+1} \to \alpha + 1$  by  $f_{\alpha+1}(x) = f_{\alpha}(x)$  for  $x \in S_{\alpha}$  and  $f_{\alpha+1}(x_{\alpha}) = \alpha$ .

Clearly,  $f_{\alpha+1}$  is an order isomorphism, and it must be unique since the maximal elements must be sent to each other and since  $f_{\alpha}$  is unique.

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<u>Limit case</u>: Let  $\lambda \leq \beta$  be a limit ordinal, and suppose that  $f_{\alpha} : S_{\alpha} \to \alpha$  is the unique order isomorphism for all  $\alpha < \lambda$ . Notice that since  $S_{\gamma} \subset S_{\delta}$  whenever  $\gamma < \delta$ , it follows that these order isomorphisms must agree. i.e., for all  $\gamma < \delta$ ,  $f_{\gamma+1}(x_{\gamma}) = f_{\delta}(x_{\gamma})$ . Now, define  $f_{\lambda} : S_{\lambda} \to \lambda$  by  $f_{\lambda}(x_{\alpha}) = f_{\alpha+1}(x_{\alpha})$ .

Suppose  $f_{\lambda}(x_{\gamma}) = f_{\lambda}(x_{\delta})$  for  $\gamma \leq \delta < \lambda$ . By definition,  $f_{\lambda}(x_{\gamma}) = f_{\gamma+1}(x_{\gamma})$  and  $f_{\lambda}(x_{\delta}) = f_{\delta+1}(x_{\delta})$ . This gives the equality

$$f_{\gamma+1}\left(x_{\gamma}\right) = f_{\delta+1}\left(x_{\delta}\right)$$

Since the functions agree on subsets,  $f_{\gamma+1}(x_{\gamma}) = f_{\delta+1}(x_{\gamma})$ . Then  $f_{\delta+1}(x_{\gamma}) = f_{\delta+1}(x_{\delta})$ , so  $x_{\gamma} = x_{\delta}$ , and  $f_{\lambda}$  is injective.

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Now, let  $\alpha \in \lambda$ . Since  $\lambda$  consists of all ordinals less than  $\lambda$ ,  $\alpha < \lambda$ . Since  $\lambda$  is a limit ordinal,  $\alpha + 1 < \lambda$ . Thus,  $f_{\lambda}(x_{\alpha}) = f_{\alpha+1}(x_{\alpha}) = \alpha$ , so  $f_{\lambda}$  is surjective.

Let  $x_{\gamma} < x_{\delta} \in S_{\lambda}$ . Then  $f_{\lambda}(x_{\gamma}) = f_{\delta+1}(x_{\gamma}) < f_{\delta+1}(x_{\delta}) = f_{\lambda}(x_{\delta})$ , so f preserves order.

Lastly,  $f_{\lambda}$  is unique by the same argument as above (since  $f_{\lambda}$  is an order isomorphism, its restriction to  $S_{\alpha}$  must be an order isomorphism with  $\alpha$ , which is unique).

Hence, by transfinite induction, there is a unique order isomorphism from  $S_{\beta}$  to  $\beta$ . Since there is clearly (by construction) a unique order isomorphism from  $(S, \leq)$  to  $S_{\beta}$ , it follows that there is a unique isomorphism from  $(S, \leq)$  to  $\beta$ , meaning that  $(S, \leq)$  has order type  $\beta$ .  $\Box$ 

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