

① Linear equations / Gaussian elimination.  
 $m, n$  integers. Solve as best possible the (linear) system

$$a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \dots = b_2$$

$$\vdots$$

$$a_{m1}x_1 + a_{m2}x_2 + \dots = b_m$$

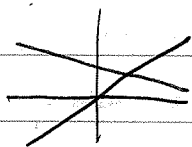
The  $\{a_{ij}\}$  and  $\{b_i\}$  are fixed. We call the system homogeneous if each  $b_i = 0$  (there is then the trivial solution) when  $n=2, 3$  and maybe 4 we set  $x_1=x, x_2=y, x_3=z, \dots$

Pictures helpful when  $n=2$  (and, less so,  $n=3$ )

$n=2$  :  $ax + a'y = b$  is a line (thru  $O$ , if  $b=0$ )

$n=3$  a plane in 3 dimensions.

Principle:  
 $n=2$



"two lines usually meet in one point". But there can be

infinitely many or none

$n=3$



"3 planes usually meet at one point."

Problems

Maybe there are really only two planes or one plane (so the solution might not exist, or could be

Ex

$$3x + 2y + 4z = 5$$

$$6x + 4y + 8z = 6$$

$$9x + 6y + 12z = 6'$$

\* But if we change the coefficients 6, 4, 8 and 9, 6, 12 just a little, there will usually be a solution

Def Singular case:  $m=n$  and 0 or only many solutions.

\*

1-2-

We can write our system as

$$x_1 \begin{bmatrix} a_{11} \\ a_{21} \\ \vdots \\ a_{m1} \end{bmatrix} + x_2 \begin{bmatrix} a_{12} \\ a_{22} \\ \vdots \\ a_{m2} \end{bmatrix} + \dots + x_n \begin{bmatrix} a_{1n} \\ a_{2n} \\ \vdots \\ a_{mn} \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_m \end{bmatrix}$$

a single linear equation.

Gaussian elimination tries to get an equivalent system (same solutions) which is easier to solve. We may

1. interchange rows
2. multiply a row by a non zero constant
3. add  $k \cdot (\text{row } s)$  to row  $t$ .

All these steps are reversible.

Goal: to have system where leading element of each row which is non zero (on the left side) is 1, with the leading elements moving to the right as we go down rows. So ( $n=4$ ) we have, say,

$$x_1 + 5x_2 + 7x_3 + 2x_4 = b_1$$

$$\nearrow x_2 - kx_3 + lx_4 = b_2$$

$$\nearrow \nearrow x_4 = b_3$$

We can see it has a solution by working up (back substitution).

Algorithm. By 1, get  $a_{11} \neq 0$ , using 2 get  $a_{11} = 1$ , then use 3 to eliminate  $a_{21}, a_{31}, \dots, a_{m1}$  (changing  $b_2, \dots, b_m$ ). Then forget first equation and continue.

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P 10 #5, 8, 9,

Principle: If  $x_1, x_2, x_3$  are 3 vectors in 3-dimensional space, then the collection

$$u x_1 + v x_2 + w x_3 \neq$$

spans 3-dimensions ( $x_1 = e_1, x_2 = e_2, x_3 = e_3$  is what you'd expect). But here

$$u \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} + v \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix} + w \begin{bmatrix} 1 \\ 3 \\ 2 \end{bmatrix}$$

doesn't work. If we take  $u=2, v=-1$ , the first two vectors already give  $\begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$ . So we get nothing extra having  $w \begin{bmatrix} 1 \\ 3 \\ 2 \end{bmatrix}$ . What we have is instead the collection of "linear combinations"

$$u \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} + v \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix},$$

so if we ask when is the system

$$u \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} + v \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix} + w \begin{bmatrix} 1 \\ 3 \\ 2 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

solvable, we work

$$u \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} + v \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix} + w (2 \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} - \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix}) = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}$$

$$\text{or } (u+2w) \begin{bmatrix} 1 \\ 1 \\ 0 \end{bmatrix} + (v-w) \begin{bmatrix} 1 \\ 2 \\ 1 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}.$$

$$\begin{bmatrix} u+v+w \\ u+2v \\ v-w \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \end{bmatrix}; \text{ we need } \boxed{b_1 = b_2 - b_3}$$

1-34

p 15 # 10

$$\begin{array}{r} 2x + 3y + z = -8 \\ 4x + 7y + 5z = 20 \\ \hline -2y - 2z = 0 \end{array}$$

$$\begin{array}{l} \textcircled{x} + \frac{3}{2}y + \frac{1}{2}z = -4 \\ (-4 \times \text{row 1}) \quad 0 + \textcircled{2y} + 3z = 36 \\ (+2 \times \text{row 2}) \quad 0 + 4z = 72 \end{array}$$

$$\begin{array}{l} \textcircled{x} + \frac{3}{2}y + \frac{1}{2}z = -4 \\ \textcircled{4} + 3z = 36 \\ \textcircled{2} = 18 \end{array}$$

How could we "fudge" this to fail?? When there is a row without a pivot.

Then there is not solution or infinitely many (identity)

Matrices, algebra.

Write original system as single matrix equation

$$Ax = b \quad \begin{matrix} \text{matrix} - \text{column vector} \\ \text{coefficient matrix} \end{matrix}$$

Basic operation: row times column vector [need same # of entries] - this is the inner product (see later)

$Ax$  by rows (simpler) or columns (pp 20-21)

$Ax$  is a sum of columns of  $A$  weighted by elements

$x$

Algebra:  $+$   $\times$  (not commutative). But associative,  $0$  matrix,  $I$ : identity matrix, distributive (check  $(EA)x = E(Ax)$ )

$E_{ij}$ : inter change

$E_i(c)$ ,  $c \neq 0$

$E_{ij}(c)$ :

$i$ th row  $\rightarrow$   $i+j$  (row)

Elementary matrices (square, depend on # of rows)

Do the operation on  $E$ , mult on left by

For step 3,  $-l$  (row  $j$ ) from row  $i$

$$E_{ij}(-l) = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -l & 0 & 1 \end{bmatrix}$$

We thus have the equation

$$E_2 \dots E_2 E_1 A = U \quad (\text{upper triangular})$$

Ex p 26 # 5  $Ax = 0$  clearly. Let's find other solutions

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 3/2 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -3 & 0 & 1 \end{bmatrix} \begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 3 & -6 & 0 \\ 0 & 2 & -2 \\ 1 & -1 & -1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 3/2 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -3 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 & -1 & -1 \\ 0 & 2 & -2 \\ 3 & -6 & 0 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 3/2 & 1 \end{bmatrix} \begin{bmatrix} 1 & -1 & -1 \\ 0 & 2 & -2 \\ 0 & -3 & 3 \end{bmatrix} \quad (\text{cont'd!})$$

2-2

$$\rightarrow \begin{bmatrix} 1 & -1 & -1 \\ 0 & 2 & -2 \\ 0 & 0 & 0 \end{bmatrix}$$

So  $Ax = 0$  becomes

$$\begin{bmatrix} 1 & -1 & -1 \\ 0 & 2 & -2 \\ 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

$$= \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}$$

$x_3$  is a free variable. Then  $x_2 = x_3$  and  $x_1 = 2x_2$

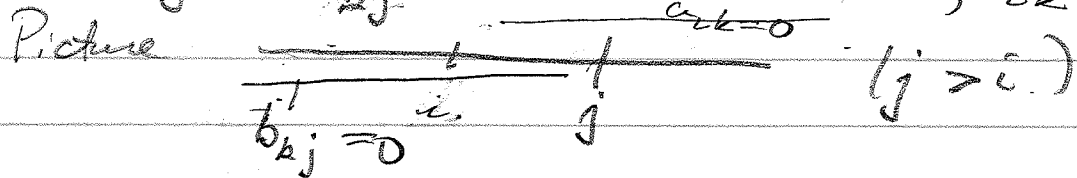
P27#12. Lower triangular means

$$a_{ij} = 0 \text{ if } j > i \text{ (check).}$$

Let  $A, B$  be lower triangular. Show  $C = AB$  is lower triangular. So let  $j > i$ , what is  $c_{ij}$ ?

$$c_{ij} = \sum_{k=1}^n a_{ik} b_{kj} = a_{i1}b_{1j} + a_{i2}b_{2j} + a_{i3}b_{3j} + \dots + a_{im}b_{mj}$$

when  $j > k$ ,  $b_{kj} = 0$  and when  $k > i$ ,  $a_{ik} = 0$



$$P27#16 \quad \underbrace{AB}_{\text{1st col of } AB} \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} = A \left( B \begin{pmatrix} 1 \\ 0 \\ \vdots \\ 0 \end{pmatrix} \right) = A \begin{pmatrix} b_{11} \\ b_{21} \\ b_{31} \\ \vdots \\ b_{m1} \end{pmatrix}$$

$A \times \text{1st col of } B$

Note  $(ABC = A(BC))$  by rarely is  $AB = BA$   
(sometimes only one will make sense!)