# Several Simple Real-world Applications of Linear Algebra Tools 

E. Ulrychová ${ }^{1}$<br>University of Economics, Department of Mathematics, Prague, Czech Republic.


#### Abstract

In this paper we provide several real-world motivated examples illustrating the power of the linear algebra tools as the product of matrices and matrix notation of systems of linear equations.


To explain some mathematical terms in a class or in a textbook it is often convenient to illustrate them by suitable examples having applications in our daily life. If the class or the textbook are on the basic level, the examples should be simple enough so that no special knowledge would be required. Four such examples, which illustrate the use of matrices, the efficiency of their products and an advantage of matrix notation for a system of linear equations, are presented.

## Example 1

Three people denoted by $P_{1}, P_{2}, P_{3}$ intend to buy some rolls, buns, cakes and bread. Each of them needs these commodities in differing amounts and can buy them in two shops $S_{1}, S_{2}$. Which shop is the best for every person $P_{1}, P_{2}, P_{3}$ to pay as little as possible? The individual prices and desired quantities of the commodities are given in the following tables:

Demanded quantity of foodstuff:

|  | roll | bun | cake | bread |
| :---: | :---: | :---: | :---: | :---: |
| $P_{1}$ | 6 | 5 | 3 | 1 |
| $P_{2}$ | 3 | 6 | 2 | 2 |
| $P_{3}$ | 3 | 4 | 3 | 1 |

Prices in shops $S_{1}$ and $S_{2}$ :

|  | $S_{1}$ | $S_{2}$ |
| ---: | ---: | ---: |
| roll | 1.50 | 1.00 |
| bun | 2.00 | 2.50 |
| cake | 5.00 | 4.50 |
| bread | 16.00 | 17.00 |

For example, the amount spent by the person $P_{1}$ in the shop $S_{1}$ is:

$$
6 \cdot 1.50+5 \cdot 2+3 \cdot 5+1 \cdot 16=50
$$

and in the shop $S_{2}$ :

$$
6 \cdot 1+5 \cdot 2.50+3 \cdot 4.50+1 \cdot 17=49
$$

for the other people similarly. These calculations can be written using a product of two matrices

$$
\mathbf{P}=\left(\begin{array}{llll}
6 & 5 & 3 & 1 \\
3 & 6 & 2 & 2 \\
3 & 4 & 3 & 1
\end{array}\right)
$$

(the demand matrix) and

$$
\mathbf{Q}=\left(\begin{array}{cc}
1.50 & 1 \\
2 & 2.50 \\
5 & 4.50 \\
16 & 17
\end{array}\right)
$$

[^0](the price matrix). For example, the first row of the matrix
\[

\mathbf{R}=\mathbf{P Q}=\left($$
\begin{array}{cc}
50 & 49 \\
58.50 & 61 \\
43.50 & 43.50
\end{array}
$$\right)
\]

expresses the amount spent by the person $P_{1}$ in the shop $S_{1}$ (the element $r_{11}$ ) and in the shop $S_{2}$ (the element $r_{12}$ ). Hence, it is optimal for the person $P_{1}$ to buy in the shop $S_{2}$, for the person $P_{2}$ in $S_{1}$ and the person $P_{3}$ will pay the same price in $S_{1}$ as in $S_{2}$.

## Example 2

To encode a short message a number can be assigned to each letter of the alphabet according to a given table. The text as a sequence of numbers will be organized into a square matrix $\mathbf{A}$; in the case that the number of letters is lower than the number of elements of the matrix $\mathbf{A}$, the rest of the matrix can be filled with zero elements. Let a nonsingular square matrix $\mathbf{C}$ be given. To encode the text the matrix $\mathbf{A}$ can be multiplied by the matrix $\mathbf{C}$ for example on the left. Let the following table and the matrix $\mathbf{C}$ be given:

| A | B | C | D | E | F | G | H | I | J | K | L | M | N | O | P | Q | R | S | T | U | V | W | X | Y | Z |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 8 | 7 | 5 | 13 | 9 | 16 | 18 | 22 | 4 | 23 | 11 | 3 | 21 | 1 | 6 | 15 | 12 | 19 | 2 | 14 | 17 | 20 | 25 | 24 | 10 | 26 |

$$
\mathbf{C}=\left(\begin{array}{lll}
2 & 0 & 1 \\
1 & 0 & 1 \\
0 & 1 & 0
\end{array}\right)
$$

We put the text "BILA KOCKA" (a white cat) into the matrix A:

$$
\mathbf{A}=\left(\begin{array}{ccc}
7 & 4 & 3 \\
8 & 11 & 6 \\
5 & 11 & 8
\end{array}\right)
$$

and encode the text:

$$
\mathbf{Z}=\mathbf{C A}=\left(\begin{array}{ccc}
19 & 19 & 14 \\
12 & 15 & 11 \\
8 & 11 & 6
\end{array}\right)
$$

To decode the message we have to multiply the matrix $\mathbf{Z}$ by the matrix $\mathbf{C}^{-\mathbf{1}}$ on the left:

$$
\mathbf{C}^{-\mathbf{1}} \mathbf{Z}=\left(\begin{array}{ccc}
1 & -1 & 0 \\
0 & 0 & 1 \\
-1 & 2 & 0
\end{array}\right)\left(\begin{array}{ccc}
19 & 19 & 14 \\
12 & 15 & 11 \\
8 & 11 & 6
\end{array}\right)=\mathbf{A}
$$

Since the matrix multiplication is not commutative, it is necessary to keep the order of the matrices in the product. If we multiply the matrices $\mathbf{C}^{\mathbf{- 1}}$ and $\mathbf{Z}$ in the opposite order, we obtain

$$
\mathbf{Z C}^{-\mathbf{1}}=\left(\begin{array}{ccc}
19 & 19 & 14 \\
12 & 15 & 11 \\
8 & 11 & 6
\end{array}\right)\left(\begin{array}{ccc}
1 & -1 & 0 \\
0 & 0 & 1 \\
-1 & 2 & 0
\end{array}\right)=\left(\begin{array}{ccc}
5 & 9 & 19 \\
1 & 10 & 15 \\
2 & 4 & 11
\end{array}\right)
$$

and it means "CERNY PSIK" (a black dog).

## ULRYCHOVÁ: SEVERAL SIMPLE REAL-WORLD APPLICATIONS OF LIN. ALG. TOOLS

## Example 3

Let us consider a group of people $P_{1}, \ldots, P_{n}$. We put $a_{i j}=1$ if the person $P_{i}$ can send some information to the person $P_{j}$, and $a_{i j}=0$ otherwise (for convenience, we put $a_{i i}=0$ for all $i=1, \ldots, n)$. Organizing this elements into a square matrix $\mathbf{A}$, we obtain a so called incidence matrix. Let $P_{i} \rightarrow P_{j}$ denotes the fact that $P_{i}$ can send information to $P_{j}$. Thus, for example, the elements of the matrix

$$
\mathbf{A}=\left(\begin{array}{llll}
0 & 1 & 0 & 1 \\
0 & 0 & 1 & 0 \\
1 & 0 & 0 & 1 \\
1 & 1 & 0 & 0
\end{array}\right)
$$

tell us that $P_{1} \rightarrow P_{2}, P_{1} \rightarrow P_{4}, P_{2} \rightarrow P_{3}, P_{3} \rightarrow P_{1}, P_{3} \rightarrow P_{4}, P_{4} \rightarrow P_{1}, P_{4} \rightarrow P_{2}$ (since $P_{1} \rightarrow P_{4}$ and $P_{4} \rightarrow P_{1}$, it is obvious that $P_{1}$ and $P_{4}$ can send information to each other). How may we interpret the matrix

$$
\mathbf{A}^{2}=\left(\begin{array}{llll}
1 & 1 & 1 & 0 \\
1 & 0 & 0 & 1 \\
1 & 2 & 0 & 1 \\
0 & 1 & 1 & 1
\end{array}\right) ?
$$

Denoting the elements of $\mathbf{A}^{2}$ by $\left(a^{2}\right)_{i j}$, we obtain for example

$$
\left(a^{2}\right)_{32}=a_{31} a_{12}+a_{32} a_{22}+a_{33} a_{32}+a_{34} a_{42}=1+0+0+1=2
$$

and this result shows that the person $P_{3}$ can send information to $P_{2}$ in two stages by two ways: $P_{3} \rightarrow P_{1} \wedge P_{1} \rightarrow P_{2}$ (because $a_{31} a_{12}=1$ ) and $P_{3} \rightarrow P_{4} \wedge P_{4} \rightarrow P_{2}$ (because $a_{34} a_{42}=1$ ). Similarly, since $\left(a^{2}\right)_{14}=0$, there is no possibility to send information from $P_{1}$ to $P_{4}$ in two stages (but it is possible directly, because $a_{14}=1$ ).

Hence, the element $\left(a^{2}\right)_{i j}$ gives the number of ways in which the person $P_{i}$ can send information to $P_{j}$ in two stages.

Similarly, $\left(a^{3}\right)_{i j}$ represents the number of ways in which the person $P_{i}$ can send information to $P_{j}$ in three stages:

$$
\mathbf{A}^{3}=\left(\begin{array}{llll}
1 & 1 & 1 & 2 \\
1 & 2 & 0 & 0 \\
1 & 2 & 2 & 1 \\
2 & 1 & 1 & 1
\end{array}\right)
$$

and thus for example

$$
\left(a^{3}\right)_{32}=\left(a^{2}\right)_{31} a_{12}+\left(a^{2}\right)_{32} a_{22}+\left(a^{2}\right)_{33} a_{32}+\left(a^{2}\right)_{34} a_{42}=1+0+0+1=2
$$

Hence, there are two ways to send information from $P_{3}$ to $P_{2}$ in three stages: $P_{3} \rightarrow P_{4} \wedge P_{4} \rightarrow$ $P_{1} \wedge P_{1} \rightarrow P_{2}$ (because $\left.\left(a^{2}\right)_{31} a_{12}=\left(a_{31} a_{11}+a_{32} a_{21}+a_{33} a_{31}+a_{34} a_{41}\right) a_{12}=(0+0+0+1) \cdot 1=1\right)$ and $P_{3} \rightarrow P_{1} \wedge P_{1} \rightarrow P_{4} \wedge P_{4} \rightarrow P_{2}$ (because $\left(a^{2}\right)_{34} a_{42}=\left(a_{31} a_{14}+a_{32} a_{24}+a_{33} a_{34}+a_{34} a_{44}\right) a_{42}=$ $(1+0+0+0) \cdot 1=1)$.

In general, the number of ways in which $P_{i}$ can send information to $P_{j}$ in at most $k$ stages is given by the element in the $i$-th row and $j$-th column of the matrix $\mathbf{A}+\mathbf{A}^{2}+\mathbf{A}^{3}+\ldots+\mathbf{A}^{k}$. Thus, in the above example we deduce from the matrix

$$
\mathbf{A}+\mathbf{A}^{2}+\mathbf{A}^{3}=\left(\begin{array}{llll}
2 & 3 & 2 & 3 \\
2 & 2 & 1 & 1 \\
3 & 4 & 2 & 3 \\
3 & 3 & 2 & 2
\end{array}\right)
$$

that for example there are four ways in which $P_{3}$ can send information to $P_{2}$ in at most three stages.

## ULRYCHOVÁ: SEVERAL SIMPLE REAL-WORLD APPLICATIONS OF LIN. ALG. TOOLS

## Example 4

Three people (denoted by $P_{1}, P_{2}, P_{3}$ ) organized in a simple closed society produce three commodities $z_{1}, z_{2}, z_{3}$. Each person sells and buys from each other. All their products are consumed by them, no other commodities enter the system (the "closed model"). The proportions of the products consumed by each of $P_{1}, P_{2}, P_{3}$ are given in the following table:

|  | $z_{1}$ | $z_{2}$ | $z_{3}$ |
| :---: | :---: | :---: | :---: |
| $P_{1}$ | 0.6 | 0.2 | 0.3 |
| $P_{2}$ | 0.1 | 0.7 | 0.2 |
| $P_{3}$ | 0.3 | 0.1 | 0.5 |

For example, the first column lists that $60 \%$ of the produced commodity $z_{1}$ are consumed by $P_{1}, 10 \%$ by $P_{2}$ and $30 \%$ by $P_{3}$. Thus, it is obvious that the sum of elements in each column is equal to 1.

Let us denote $x_{1}, x_{2}, x_{3}$ the incomes of the persons $P_{1}, P_{2}, P_{3}$. Then the amount spent by $P_{1}$ on $z_{1}, z_{2}, z_{3}$ is $0.6 x_{1}+0.2 x_{2}+0.3 x_{3}$. The assumption that the consumption of each person equals his income leads to the equation $0.6 x_{1}+0.2 x_{2}+0.3 x_{3}=x_{1}$, similarly for the other persons. We obtain the system of linear equations:

$$
\begin{aligned}
& 0.6 x_{1}+0.2 x_{2}+0.3 x_{3}=x_{1} \\
& 0.1 x_{1}+0.7 x_{2}+0.2 x_{3}=x_{2} \\
& 0.3 x_{1}+0.1 x_{2}+0.5 x_{3}=x_{3} .
\end{aligned}
$$

This system can be rewritten as the equation $\mathbf{A x}=\mathbf{x}$, where

$$
\mathbf{A}=\left(\begin{array}{ccc}
0.6 & 0.2 & 0.3 \\
0.1 & 0.7 & 0.2 \\
0.3 & 0.1 & 0.5
\end{array}\right) \quad \text { and } \quad \mathbf{x}=\left(x_{1}, x_{2}, x_{3}\right)^{T}
$$

Moreover, we assume the income to be nonnegative, i.e. $x_{i} \geq 0$ for $i=1,2,3$ (we denote it $\mathbf{x} \geq \mathbf{o})$. We can rewrite this equation into the equivalent form $(\mathbf{A}-\mathbf{I}) \mathbf{x}=\mathbf{o}$ :

$$
\left(\begin{array}{ccc|c}
-0.4 & 0.2 & 0.3 & 0 \\
0.1 & -0.3 & 0.2 & 0 \\
0.3 & 0.1 & -0.5 & 0
\end{array}\right)
$$

An arbitrary solution of the system has the form $\mathbf{x}=t(13,11,10)^{T}$ and it is $\mathbf{x} \geq \mathbf{o}$ for $t \geq 0$.
Thus, to ensure that this society survives, the persons $P_{1}, P_{2}, P_{3}$ have to have their incomes in the proportions 13:11:10.

Note: Let us consider a closed model, let $\mathbf{A}=\left(a_{i j}\right)$ be an $n \times n$ coefficient matrix as above. Instead of the condition that the consumption is equal to the income we can consider that the consumption does not exceed the income, i.e. $a_{i 1} x_{1}+a_{i 2} x_{2}+\ldots+a_{i n} x_{n} \leq x_{i}$ for all $i=1, \ldots, n$. We will show that the equation $\mathbf{A x}=\mathbf{x}$ has to hold in this case, too. For otherwise, there exists a $p$ such that $a_{p 1} x_{1}+a_{p 2} x_{2}+\ldots+a_{p n} x_{n}<x_{p}$ and then $\sum_{i=1}^{n}\left(a_{i 1} x_{1}+\ldots a_{i n} x_{n}\right)<\sum_{i=1}^{n} x_{i}$. Since the sum of elements in each column of the matrix $\mathbf{A}$ is equal to 1 , the left side of the previous unequality can be rewritten: $\sum_{i=1}^{n}\left(a_{i 1} x_{1}+\ldots+a_{i n} x_{n}\right)=\sum_{i=1}^{n} \sum_{j=1}^{n} a_{i j} x_{j}=\sum_{j=1}^{n}\left(\sum_{i=1}^{n} a_{i j}\right) x_{j}=$ $\sum_{j=1}^{n} 1 \cdot x_{j}=\sum_{j=1}^{n} x_{j}$. Thus $\sum_{j=1}^{n} x_{j}<\sum_{i=1}^{n} x_{i}$ and it is a contradiction.

## References

Friedberg, S.H., Insel, A.J., Spence, L.E.: Linear Algebra, 4th edition, Prentice Hall, 2003.
Poole, D.: Linear Algebra: A Modern Introduction. Brooks/Cole, 2003.
Coufal, J., Ulrychová, E.: Šifrování a dešifrování užitím regulárních matic, Mundus Symbolicus, 4, 15-19, 1996.


[^0]:    ${ }^{1}$ email: uleva@vse.cz
    The author is a postgradual student at Charles University, Faculty of Mathematics and Physics, Prague, Czech Republic.

