

Meeting 10.2: Topological quantum computing, I

I. Anyons from elementary excitations in toric code

II. TQC and TQFT go hand in hand.

I. Anyons from elementary excitations in toric code

Given a cellulation w/ N edges of g genus g surface S_g , toric code yields a 4^g -dimensional code space

$$\mathcal{H} \subseteq (\mathbb{C}^2)^{\otimes N}$$

$$\mathcal{H} = \left\{ |\psi\rangle \mid X_v |\psi\rangle = |\psi\rangle = Z_p |\psi\rangle \text{ for vertices } v, \text{ plaquettes } p \right\}$$

We can repackage \mathcal{H} as the "ground state space" of

$$H = \sum_v (I - X_v) + \sum_p (I - Z_p)$$

$$\text{i.e. } \mathcal{H} = \ker H.$$

\nwarrow H is a nonnegative Hermitian operator

$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$\text{spec } X = \{-1, 1\}$$

$$I - X = \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$$

$$\text{spec } I - X = \{0, 2\}$$

Similarly,

$$\text{spec } I - X_p = \{0, 2\}$$

$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

$$\text{spec } Z = \{-1, 1\}$$

$$I - Z = \begin{pmatrix} 0 & 0 \\ 0 & 2 \end{pmatrix}$$

$$\text{spec } I - Z = \{0, 2\}$$

$$\text{spec } I - Z_p = \{0, 2\}$$

X_V 's and Z_p 's commute, hence so do

$I - X_V$'s and $I - Z_p$'s

Physics intuition: \mathcal{H} consists of "vacuum" or "zero-energy" states w.r.t. Hamiltonian H .

Eigenvectors of H corresponding to Non-zero eigenvalues?

$$\text{spec } H: 0 < \lambda_1 < \lambda_2 < \dots < \lambda_k$$

\uparrow

The eigenspace is the state

space for (pairs of) "elementary particles" of this system...

$$\text{In fact, spec } H: 0 < 4 < 8 < \dots$$

What is E_{λ_1} , the eigenspace corresponding to λ_1 ?

More intuition: suppose $X_v |q\rangle \neq |q\rangle$.

We might say that $|q\rangle$ has a "charge" or "particle" at vertex v .
(nonzero)

Similarly, if $\sum_P |q\rangle \neq |q\rangle$, might say $|q\rangle$ has a (nonzero) flux through P , or $|q\rangle$ has a "vortex" on plaquette P .

A lowest energy state, i.e. $|q\rangle \in E_i$, must violate as few of the constraints $X_v |q\rangle = |q\rangle$, $\sum_P |q\rangle = |q\rangle$ as possible.

B/c $\prod_v X_v = I = \prod_P \sum_P |q\rangle$ either violates exactly two X_v 's or two \sum_P 's.

Recall:

If c is a loop in 1-skeleton,

define

$$\mathcal{Z}_c = \prod_{e \in c} \mathcal{Z}_e$$

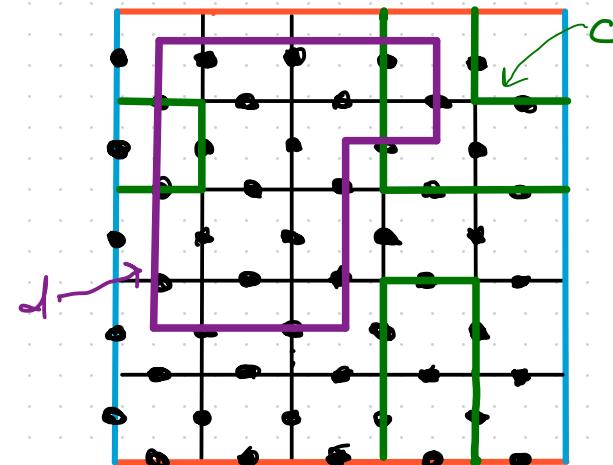
If d a loop in dual 1-skeleton,

define

$$X_d = \prod_{e \in d} X_e$$

The loop operators generate possible errors of the code.

A loop operator implements an undetectable (and non-trivial) error if and only if the loop is non-trivial in $H_1(S^1 \times S^1; \mathbb{Z}/2)$



String Operators:

If c is a path in 1-skeleton,
define

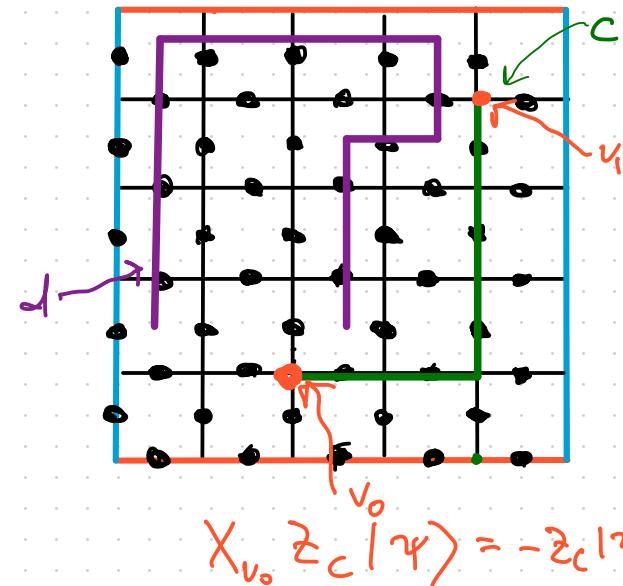
$$\mathcal{Z}_c = \prod_{e \in c} \mathcal{Z}_e$$

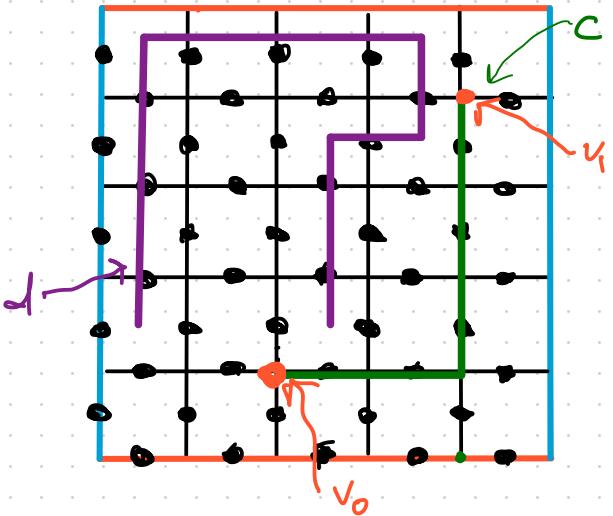
If d a path in dual 1-skeleton,
define

$$X_d = \prod_{e \in d} X_e.$$

If $|v\rangle \in \mathcal{H}$, then $\mathcal{Z}_c |v\rangle$ and $X_d |v\rangle$ are in

E_{λ_1} . For instance, $\mathcal{Z}_c |v\rangle$ will violate the two X_v 's at ends of c .





$$X_{v_0} z_c | \Psi \rangle = -z_c | \Psi \rangle$$

$$X_{v_1} z_c | \Psi \rangle = -z_c | \Psi \rangle$$

On the other hand

$$X_v z_c | \Psi \rangle = z_c | \Psi \rangle$$

for any $v \neq v_0, v_1$

and

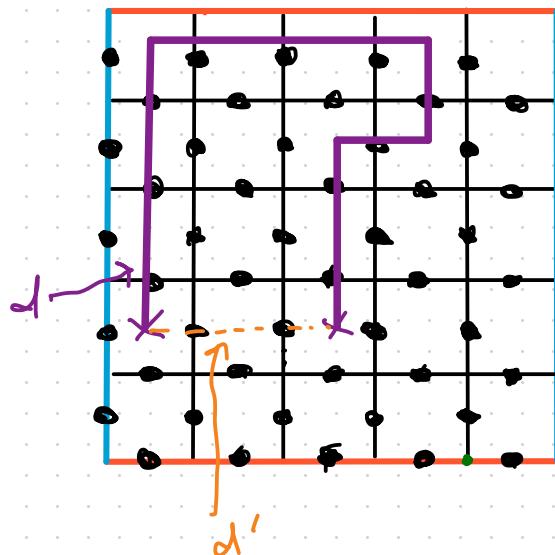
$$Z_P z_c | \Psi \rangle = z_c | \Psi \rangle$$

for all P .

Note: If $|q\rangle \in \mathcal{H}_c$

$\mathcal{Z}_c |q\rangle = \mathcal{Z}_{c'} |q\rangle$ if and only if
 $c + c' = 0$ in $H_1(S^1 \times S^1; \mathbb{Z}/2)$

Similarly for X_α 's.



More generally,
we need
 $c \sim c'$,
rel
endpoints
i.e., for other
surface codes?

Kitarev's (\$100,000,000?) idea:

Introduce a small number of particles onto the surface, and move them around in controlled ways in order to intentionally manipulate a codestate of the toric code. Because nontrivial operations occur only after doing something "topologically nontrivial," the probability of implementing the wrong operation can be made small without much overhead.

In other words: "Fault tolerance from hardware" if you can implement the toric code Hamiltonian in a lab.

Two related ideas for how to process
information topologically using toric code:

1. braiding

2. "Dehn twisting"

Braiding in Kitaev's model) $|q\rangle \in \mathcal{H}$

Start w/ state:

$$X_d |z_c| q\rangle$$

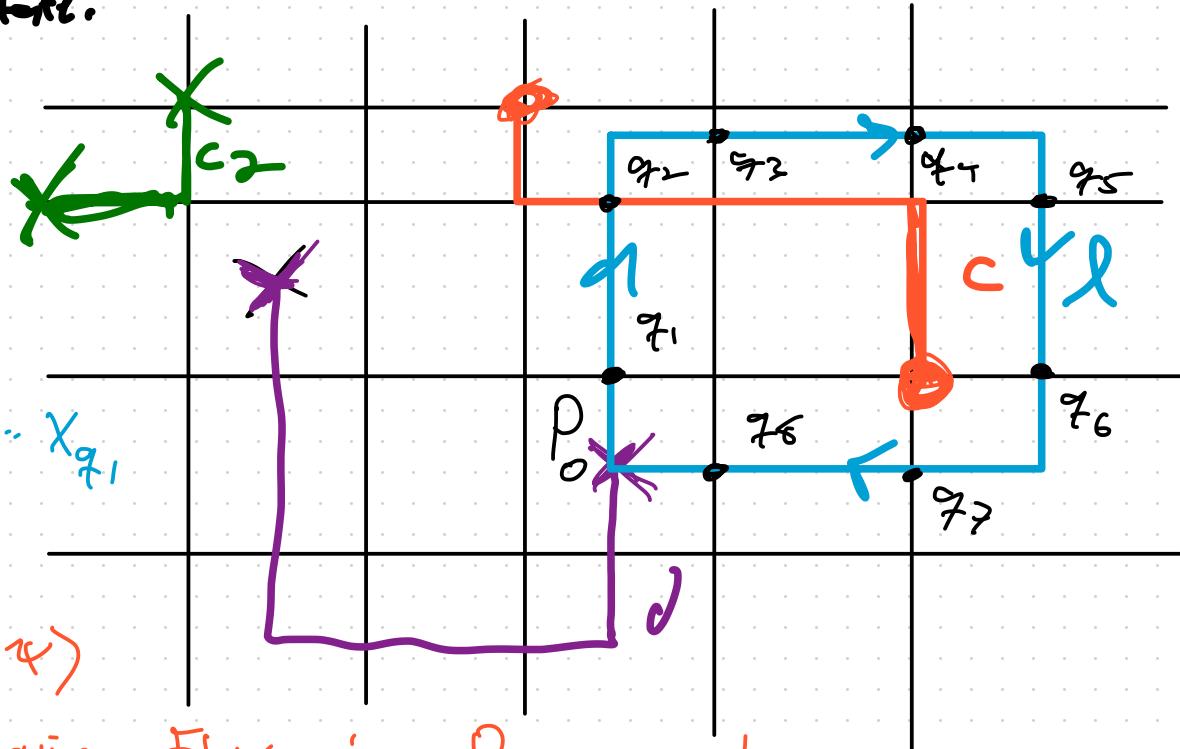
$$E_{x_2}$$

$$X_l = X_{q_0} X_{q_1} \dots X_{q_l}$$

Applying X_l
to $X_d |z_c| q\rangle$

is like moving Flux in P_0 ground

one of the charges at the end of the string



$$X_\ell(X_j z_c | \psi \rangle) = (-1)^j X_j z_c | \psi \rangle$$

Note: this is true only because
of the charge at the end of
the string c.

Weird! Moving Flux along a
loop around a charge implements
a nontrivial charge to the state.

The elementary violations of
the Stabilizer (i.e. the
particles) have "non trivial"
braiding statistics.

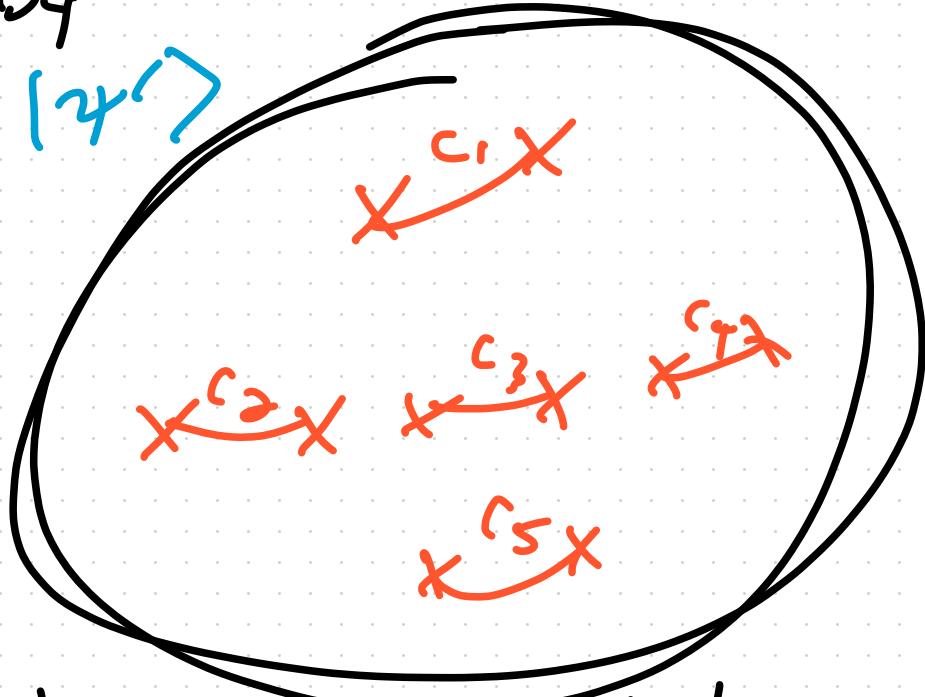
Such particles are called
anyons.

Start w) $|4\rangle \in \mathcal{H}$.

"Create particle" by

$$\sum_{c_1, c_2, \dots, c_5} |4\rangle = |\overline{4}\rangle$$

This, it does 10
particles and puts
us inside E_{15}



Loop operators yield a representation
of 10 strand braid group!

$\langle \gamma \rangle \in E_{\lambda_5}$

$V = \text{Loops} \cdot \langle \gamma \rangle \subseteq E_{\lambda_5}$

Unitary
Representations

$B_p \longrightarrow U(V)$.

"Dehn twisting" $|00\rangle \mapsto \mathbb{Z}_{C_0} |00\rangle \mapsto \mathbb{Z}_c, \mathbb{Z}_{C_0} |00\rangle$

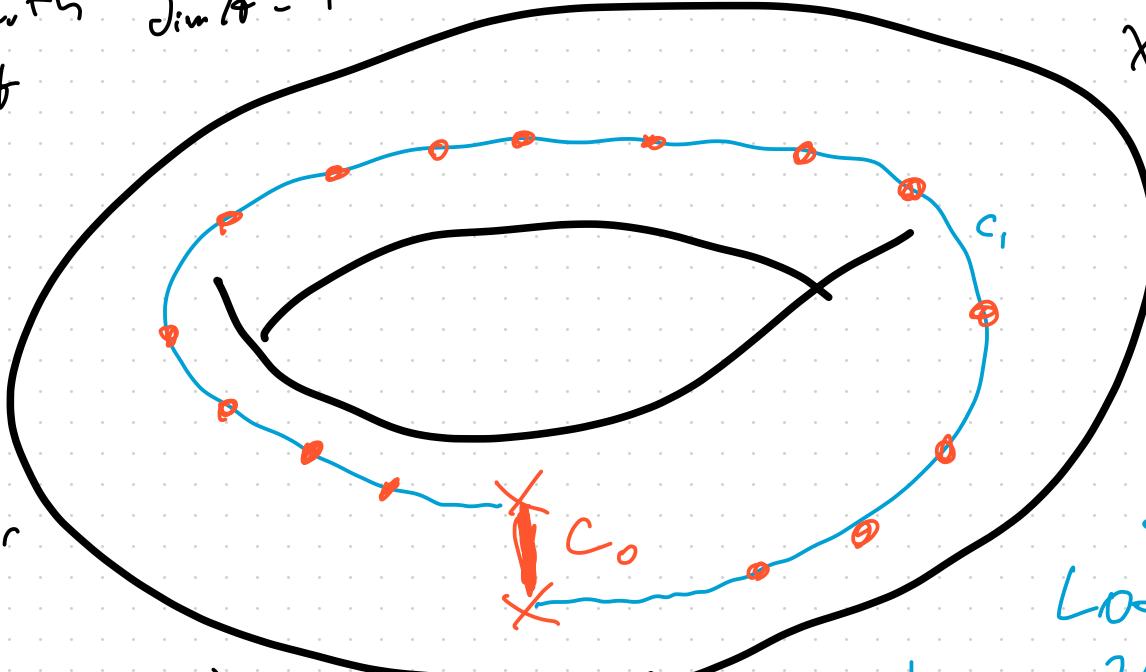
Start with $\dim \mathcal{H} = 4$

$|11\rangle \in \mathcal{H}$

Manipulate
 $|11\rangle$ into
another state
in \mathcal{H} by
creating pair
of anyons,

move one around

for us, then annihilate



Recall:

Loop operators
act on \mathcal{H} like
the Pauli X 's
and Σ 's

$$\mathcal{H} \cong \mathbb{C}^2 \otimes \mathbb{C}^2 = \text{Span} \left\{ |00\rangle, |01\rangle, |10\rangle, |11\rangle \right\}$$

loop operators act on \mathcal{H} by

$$X \otimes I_2, Z \otimes I_2$$

$$I_2 \otimes X, \text{ or } I_2 \otimes Z$$

So, can process quantum info in \mathcal{H}
by applying loop operators.

Problems with toric code?

With toric code, we can only implement X's and Z's

on codespace using loop operators. So "Dehus trust"

idea is insufficient to generate quantum universal op's on it

Similar issues for bridging

Are there other, similar setups,
but where the topologically protected
operations are powerful enough to
implement a quantum universal
gate set?

Yes!

This is what Freedman - Larsen - Wang
prove.

II. TQC and TQFT go hand in hand.

Once-extended $(2+1)$ -dimensional
TQFTs provide a language to
abstract away the combinatorial
aspects of Kitaev's proposal,
and focus on the topology
of anyons and their interactions