

# Solovay-Kitaev Theorem (SU(2) version)

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$\mathcal{G}$  is a finite gate set w/ inverses,  $\langle \mathcal{G} \rangle$  dense in  $SU(2)$

Thm: Any  $U \in SU(2)$  can be approx'ed within dist  $\epsilon$

using  $O(\text{polylog}(\frac{1}{\epsilon}))$  gates from  $\mathcal{G}$ .

$O(\text{polylog}(\frac{1}{\epsilon}))$  gates from  $\mathcal{G}$ .  
 $\text{polylog}(\frac{1}{\epsilon})^c$  for some const  $c$

Why care? Exists such  $\mathcal{G}$  that is fault-tolerant

$\Rightarrow$  any quantum circuit w/  $f(n)$  gates can be approx'ed within dist  $\epsilon$  w/ fault-tolerance w/  $O(f(n)\text{polylog}(f(n)/\epsilon))$  gates.

Notation:  $\mathcal{G}_l =$  length  $l$  words from  $\mathcal{G}$

$$(\langle \mathcal{G} \rangle) = \bigcup_{l \geq 0} \mathcal{G}_l$$

$S$   $\epsilon$ -net of  $T$  means  $\forall x \in T \exists y \in S$  at most  $\epsilon$  distance away

$B_\epsilon$  ball of radius  $\epsilon$  around  $I$ .

$$\text{distance tr norm } D(u, v) = \|u - v\|$$

$$\|x\| = \sqrt{x^T x}$$

Step 1(a) approx  $I$  as well as desired.

depends on  $SU(2)$  [Key Lemma: For sufficiently small  $\epsilon$ , if  $\mathcal{G}_l$  is  $\epsilon^2$ -net for  $B_\epsilon$ , then  $\mathcal{G}_{3l}$  is  $C\epsilon^3$ -net for  $B_{C\epsilon^3}$ ]

Since  $\langle \mathcal{G} \rangle$  dense  $\exists l_0, \epsilon, \epsilon_0$  s.t.  $\mathcal{G}_{l_0}$   $\epsilon_0^2$ -net for  $B_{\epsilon_0}$ .

By induction  $\mathcal{G}_{5^{k+1}l_0}$   $\epsilon_k^2$ -net for  $B_{\epsilon_k}$

$$\epsilon_k = \frac{(C\epsilon_0)^{(3/2)^k}}{C} \quad \text{wlog } \epsilon_k^2 < \epsilon_{k+1}.$$

Step 1(b) approx arbitrary  $U$  as well as desired.

$\langle \mathcal{G} \rangle$  dense so  $U_0$   $\epsilon_0^2$ -approx of  $U$

By induction:  $U_k U_{k-1} \dots U_0$   $\epsilon_k^2$ -approx of  $U$

$$U_i \in \mathcal{G}_{5^{i+1}l_0} \dots$$

$$U_i \in \mathcal{G}_{S^i l_0} \quad U = U_i U_{i-1} \dots U_0$$

$$U = U_i U_{i-1} \dots U_0 = U_i U_0^\dagger U_1^\dagger \dots U_{i-1}^\dagger$$

$V$  is within dist  $\epsilon_k^2$  of  $I$

$$V \in B_{\epsilon_{k+1}} \quad \exists U_{k+1} \in \mathcal{G}_{S^{k+1} l_0} \quad \epsilon_{k+1}^2 - \text{approx of } V.$$

$U_{k+1} U_k \dots U_0$   $\epsilon_{k+1}^2$ -approx of  $U$ .

Total # of gates for  $\epsilon_k^2$ -approx is  $\sum_{i=0}^k S^i l_0 = \frac{S^{k+1}}{4} l_0 = O(S^k)$

For desired approx  $\epsilon$ , find  $\epsilon_k^2 < \epsilon$

$$\frac{(C \epsilon_0)^{(3/2)^k}}{C} < \epsilon$$

const

$$\left(\frac{3}{2}\right)^k \frac{1}{C} \log\left(\frac{1}{\epsilon_0}\right) < \log\left(\frac{1}{\epsilon}\right)$$

$$\left(\frac{3}{2}\right)^k = O(\log(\frac{1}{\epsilon}))$$

$$\log_{3/2} 5 \approx 4$$

$$\Rightarrow S^k = O(\log(\frac{1}{\epsilon})^4).$$

Step 0.

Key Lemma: For sufficiently small  $\epsilon$ , if  $\mathcal{G}_\epsilon$  is  $\epsilon^2$ -net for  $B_\epsilon$ , then  $\mathcal{G}_{4\epsilon}$  is  $C\epsilon^3$ -net for  $B_{4\epsilon}$

Step 0(a) If  $\mathcal{G}_\epsilon$  is  $\epsilon^2$ -net for  $B_\epsilon$   
then  $\mathcal{G}_{4\epsilon}$  is  $\epsilon^2$ -net for  $B_{4\epsilon}$

Use Fact about  $SU(2)$ :

$$(1) \quad \text{If } U \in SU(2) \quad U = u(\vec{a}) := \exp(-i \vec{a} \cdot \vec{\sigma}/2)$$

$$(2) \quad [\vec{a} \cdot \vec{\sigma}, \vec{b} \cdot \vec{\sigma}] = i[\vec{a}, \vec{b}] = i\vec{a} \cdot \vec{b} - \vec{b} \cdot \vec{a}$$

Real vector  
Even part indices

$$(2) [a \cdot \sigma, b \cdot \sigma] = (\underbrace{[a, b]}_{= AB - BA}) \cdot \sigma$$

$$= 2i(\vec{a} \times \vec{b}) \cdot \vec{\sigma}$$

$\vec{a}, \vec{b}$  vector  
 $\vec{\sigma}$  Pauli matrix  
 $\vec{a} \cdot \vec{\sigma}$  is a matrix

If  $U$  is close to  $I$  then

(3i)  $U = \exp(-iA)$   $A$  hermitian w/ very small trace.

(3ii)  $\exp(-iA)\exp(-iB)\exp(iA)\exp(iB) \approx \exp(-[A, B])$

Group commutator

$O(\varepsilon^0)$  approx

Taylor expand + small traces

Prob  $U \in B_{\varepsilon^2}$ .  $U = u(\vec{a})$   $\|\vec{a}\| < \varepsilon^2$

Find Two vectors  $\vec{b} \neq \vec{c}$  s.t.  $\vec{a} = \vec{b} \times \vec{c}$   $\|\vec{b}\|, \|\vec{c}\| < \varepsilon$

$\rightarrow u(\vec{b}) u(\vec{c}) \in B_\varepsilon$

Find  $\underbrace{u(\vec{x}), u(\vec{y})}_{\in \mathcal{L}_\varepsilon}$   $\varepsilon^2$ -approx for  $u(\vec{b}), u(\vec{c})$

$\rightarrow u(\vec{x} \times \vec{y}) O(\varepsilon^0)$ -approx for  $u(\vec{b} \times \vec{c}) = u(\vec{a})$ .

$$U = u(\vec{a}) \approx u(\vec{x} \times \vec{y}) = \exp\left(2i(\vec{x} \times \vec{y}) \cdot \vec{\sigma}/4\right)$$

$$\stackrel{(2)}{=} \exp\left(-\left[\frac{\vec{x} \cdot \vec{\sigma}}{2}, \frac{\vec{y} \cdot \vec{\sigma}}{2}\right]\right)^{[\vec{x} \cdot \vec{\sigma}, \vec{y} \cdot \vec{\sigma}]}$$

$$\stackrel{(3ii)}{\approx} \exp(-i\vec{x} \cdot \vec{\sigma}/2) \dots$$

$$O(\varepsilon^3) \text{-approx} = u(\vec{x}) u(\vec{y}) u(\vec{x})^+ u(\vec{y})^+$$

$$\in \mathcal{L}_{4\varepsilon}$$

Step 0(b)  $\rightarrow \mathcal{L}_{5\varepsilon}$   $O(\varepsilon^3)$ -net for  $B_{O(\varepsilon^{3/2})}$ .

$U \in B_{O(\varepsilon^{3/2})} \subseteq B_\varepsilon \vee \in \mathcal{L}_\varepsilon$   $\varepsilon^2$ -approx of  $U$ .

$$UV^{-1} = UV^t \in B_{\varepsilon^2} \quad W \in \mathcal{G}_{4\ell} \quad O(\varepsilon^3) - \text{approx for } UV^t$$

$$WV \quad O(\varepsilon^3) - \text{approx for } U$$

$$\in \mathcal{D}_{4\ell}, \mathcal{G}_\ell = \mathcal{G}_{5\ell}$$

□