

Combinatorial Complexity in O-minimal Geometry

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Outline

- 1 Introduction
 - Some basic results
 - Combinatorial and Algebraic Complexity
- 2 Arrangements
- 3 O-minimal Structures and Admissible Sets
 - Examples of Admissible Sets
 - \mathcal{A} -sets
- 4 Results
 - Bounds on Betti Numbers
 - Cylindrical Definable Cell Decomposition
 - Application: Generalization of a Theorem due to Alon et al.
- 5 Idea of Proofs

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Semi-algebraic Sets and their Betti numbers

- Let $S \subset \mathbb{R}^k$ be defined by a Boolean formula whose atoms consists of $P > 0, P = 0, P < 0, P \in \mathcal{P}$, where \mathcal{P} is a set of polynomials of degrees bounded by a parameter and $\#\mathcal{P} = n$.



$$\sum_{i \geq 0} b_i(S) \leq n^{2k} O(d)^k. \quad (\text{Gabrielov-Vorobjov, 2005})$$

- Bound for sign conditions: (B-Pollack-Roy, 2005)

$$\sum_{\sigma \in \{0,1,-1\}^{\mathcal{P}}} b_i(\mathcal{R}(\sigma)) \leq \sum_{j=0}^{k-i} \binom{n}{j} 4^j d (2d-1)^{k-1} = n^{k-i} O(d)^k.$$

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Combinatorial Complexity

- Notice that the bounds in the previous page are products of two quantities – one that depends only on n (and k), and another part which is independent of n . We refer to the first part as the **combinatorial part** of the complexity, and the latter as the **algebraic part**.
- While understanding the **algebraic part** of the complexity is a very important problem, in several applications, most notably in **discrete and computational geometry**, it is the **combinatorial part** of the complexity that is of interest (the algebraic part is assumed to be bounded by a constant).

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Definition of Arrangements

- Let $\mathcal{A} = \{S_1, \dots, S_n\}$, with each S_i belonging to some “simple” class of sets.
- For $I \subset \{1, \dots, n\}$, let $\mathcal{A}(I)$ denote the set

$$\bigcap_{i \in I} S_i \cap \bigcap_{j \in [1 \dots n] \setminus I} \mathbb{R}^k \setminus S_j,$$

and it is customary to call a connected component of \mathcal{A}_I a **cell** of the arrangement \mathcal{A} and we denote by $\mathcal{C}(\mathcal{A})$ the set of all non-empty cells of the arrangement \mathcal{A} .

- The cardinality of $\mathcal{C}(\mathcal{A})$ is called the **combinatorial complexity** of the arrangement \mathcal{A} .

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Objects of Bounded Description Complexity

- The class of sets usually considered in the study of arrangements are sets with “**bounded description complexity**”. This means that each set in the arrangement is defined by a first order formula in the language of ordered fields involving at most a constant number polynomials whose degrees are also bounded by a constant.
- Additionally, there is often a requirement that the sets be in “**general position**”. The precise definition of “general position” varies with context, but often involves restrictions such as: the sets in the arrangements are smooth manifolds, intersecting transversally.

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Definition of O-minimal Structures

An o-minimal structure on a real closed field \mathbf{R} is a sequence $\mathcal{S}(\mathbf{R}) = (\mathcal{S}_n)_{n \in \mathbb{N}}$.

- 1 All algebraic subsets of \mathbf{R}^n are in \mathcal{S}_n .
- 2 The class \mathcal{S}_n is closed under complementation and finite unions and intersections.
- 3 If $A \in \mathcal{S}_m$ and $B \in \mathcal{S}_n$ then $A \times B \in \mathcal{S}_{m+n}$.
- 4 If $\pi : \mathbf{R}^{n+1} \rightarrow \mathbf{R}^n$ is the projection map on the first n co-ordinates and $A \in \mathcal{S}_{n+1}$, then $\pi(A) \in \mathcal{S}_n$.
- 5 The elements of \mathcal{S}_1 are precisely finite unions of points and intervals.

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Examples of O-minimal Structures I

- Our first example of an o-minimal structure $\mathcal{S}(\mathbb{R})$, is the o-minimal structure over a real closed field \mathbb{R} where each \mathcal{S}_n is exactly the class of semi-algebraic subsets of \mathbb{R}^n .
- Let \mathcal{S}_n be the images in \mathbb{R}^n under the projection maps $\mathbb{R}^{n+k} \rightarrow \mathbb{R}^n$ of sets of the form $\{(\mathbf{x}, \mathbf{y}) \in \mathbb{R}^{n+k} \mid P(\mathbf{x}, \mathbf{y}, \mathbf{e}^{\mathbf{x}}, \mathbf{e}^{\mathbf{y}}) = 0\}$, where P is a real polynomial in $2(n+k)$ variables, and $\mathbf{e}^{\mathbf{x}} = (e^{x_1}, \dots, e^{x_n})$ and $\mathbf{e}^{\mathbf{y}} = (e^{y_1}, \dots, e^{y_k})$. We will denote this o-minimal structure over \mathbb{R} by $\mathcal{S}_{\text{exp}}(\mathbb{R})$.

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- Let \mathcal{S}_n be the images in \mathbb{R}^n under the projection maps $\mathbb{R}^{n+k} \rightarrow \mathbb{R}^n$ of sets of the form

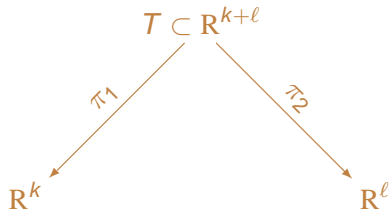
$\{(\mathbf{x}, \mathbf{y}) \in \mathbb{R}^{n+k} \mid P(\mathbf{x}, \mathbf{y}) = 0\}$, where P is a **restricted analytic function** in $2(n+k)$ variables.

(A restricted analytic function in N variables is an analytic function defined on an open neighborhood of $[0, 1]^N$ restricted to $[0, 1]^N$ (and extended by 0 outside)).

We will denote this o-minimal structure over \mathbb{R} by $\mathcal{S}_{\text{ana}}(\mathbb{R})$.

Admissible Sets

- Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure on a real closed field \mathbb{R} and let $T \subset \mathbb{R}^{k+l}$ be a fixed definable set.



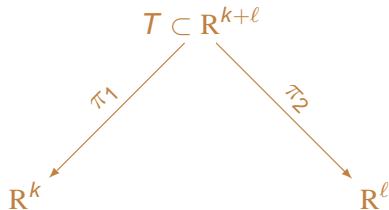
- We will call S of \mathbb{R}^k to be a (T, π_1, π_2) -set if

$$S = T_{\mathbf{y}} = \pi_1(\pi_2^{-1}(\mathbf{y}) \cap T)$$

for some $\mathbf{y} \in \mathbb{R}^l$.

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Example I

Let $\mathcal{S}(\mathbb{R}) = \mathcal{S}_{\text{sa}}(\mathbb{R})$ and Let $T \subset \mathbb{R}^{2k+1}$ be the semi-algebraic set defined by

$$T = \{(x_1, \dots, x_k, a_1, \dots, a_k, b) \mid \langle \mathbf{a}, \mathbf{x} \rangle - b = 0\}$$

(where we denote $\mathbf{a} = (a_1, \dots, a_k)$ and $\mathbf{x} = (x_1, \dots, x_k)$), and π_1 and π_2 are the projections onto the first k and last $k+1$ co-ordinates respectively. A (T, π_1, π_2) -set is clearly a hyperplane in \mathbb{R}^k and vice versa.

Example II

Let $\mathcal{S}(\mathbb{R}) = \mathcal{S}_{\text{exp}}(\mathbb{R})$ and

$$T = \{(\mathbf{x}, \mathbf{y}_1, \dots, \mathbf{y}_m, \mathbf{a}_1, \dots, \mathbf{a}_m) \mid \mathbf{x}, \mathbf{y}_1, \dots, \mathbf{y}_m \in \mathbb{R}^k, \\ \mathbf{a}_1, \dots, \mathbf{a}_m \in \mathbb{R}, x_1, \dots, x_k > 0, \sum_{i=0}^m a_i \mathbf{x}^{y_i} = 0\},$$

with $\pi_1 : \mathbb{R}^{k+m(k+1)} \rightarrow \mathbb{R}^k$ and $\pi_2 : \mathbb{R}^{k+m(k+1)} \rightarrow \mathbb{R}^{m(k+1)}$ be the projections onto the first k and the last $m(k+1)$ co-ordinates respectively. The (T, π_1, π_2) -sets in this example include (amongst others) all semi-algebraic sets consisting of intersections with the positive orthant of all real algebraic sets defined by a polynomial having at most m monomials (different sets of monomials are allowed to occur in different polynomials).

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\mathcal{A} -sets I

Let $\mathcal{A} = \{S_1, \dots, S_n\}$, such that each $S_i \subset \mathbb{R}^k$ is a (T, π_1, π_2) -set. For $I \subset \{1, \dots, n\}$, we let $\mathcal{A}(I)$ denote the set

$$\bigcap_{i \in I} S_i \cap \bigcap_{j \in [1 \dots n] \setminus I} \mathbb{R}^k \setminus S_j, \quad (1)$$

and we will call such a set to be a **basic \mathcal{A} -set**. We will denote by $\mathcal{C}(\mathcal{A})$, the set of non-empty connected components of all basic \mathcal{A} -sets.

\mathcal{A} -sets II

We will call definable subsets $S \subset \mathbb{R}^k$ defined by a Boolean formula whose atoms are of the form, $x \in S_i, 1 \leq i \leq n$, a \mathcal{A} -set. A \mathcal{A} -set is thus a union of basic \mathcal{A} -sets. If T is closed, and the Boolean formula defining S has no negations, then S is closed by definition (since each S_i is closed) and we call such a set an \mathcal{A} -closed set.

Moreover, if V is any closed definable subset of \mathbb{R}^k , and S is an \mathcal{A} -set (resp. \mathcal{A} -closed set), then we will call $S \cap V$ to be an (\mathcal{A}, V) -set (resp. (\mathcal{A}, V) -closed set).

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Bounds on Betti Numbers I

Theorem

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure over a real closed field \mathbb{R} and let $T \subset \mathbb{R}^{k+\ell}$ be a closed definable set. Then, there exists a constant $C = C(T) > 0$ depending only on T , such that for any (T, π_1, π_2) -family $\mathcal{A} = \{S_1, \dots, S_n\}$ of subsets of \mathbb{R}^k the following holds. For every $i, 0 \leq i \leq k$,

$$\sum_{D \in \mathcal{C}(\mathcal{A})} b_i(D) \leq C \cdot n^{k-i}.$$

In particular, the combinatorial complexity of \mathcal{A} , is at most $C \cdot n^k$. The topological complexity of any m cells in the arrangement \mathcal{A} is bounded by $m + C \cdot n^{k-1}$.

Lower dimensional

Theorem

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure over a real closed field \mathbb{R} and let $T \subset \mathbb{R}^{k+\ell}$, $V \subset \mathbb{R}^k$ be closed definable sets with $\dim(V) = k'$. Then, there exists a constant $C = C(T, V) > 0$ depending only on T and V , such that for any (T, π_1, π_2) -family, $\mathcal{A} = \{S_1, \dots, S_n\}$, of subsets of \mathbb{R}^k , and for every i , $0 \leq i \leq k'$,

$$\sum_{D \in \mathcal{C}(\mathcal{A}, V)} b_i(D) \leq C \cdot n^{k'-i}.$$

In particular, the combinatorial complexity of \mathcal{A} restricted to V , is bounded by $C \cdot n^{k'}$.

Topological Complexity of \mathcal{A} -sets

Theorem

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure over a real closed field \mathbb{R} , and let $T \subset \mathbb{R}^{k+\ell}$, $V \subset \mathbb{R}^k$ be closed definable sets with $\dim(V) = k'$. Then, there exists a constant $C = C(T, V) > 0$ such that for any (T, π_1, π_2) -family, \mathcal{A} with $|\mathcal{A}| = n$, and an \mathcal{A} -closed set $S_1 \subset \mathbb{R}^k$, and an \mathcal{A} -set $S_2 \subset \mathbb{R}^k$,

$$\sum_{i=0}^{k'} b_i(S_1 \cap V) \leq C \cdot n^{k'} \text{ and,}$$

$$\sum_{i=0}^{k'} b_i(S_2 \cap V) \leq C \cdot n^{2k'}.$$

Topological Complexity of Projections

Theorem (Topological Complexity of Projections)

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure, and let $T \subset \mathbb{R}^{k+\ell}$ be a definable, closed and bounded set. Let $k = k_1 + k_2$ and let $\pi_3 : \mathbb{R}^k \rightarrow \mathbb{R}^{k_2}$ denote the projection map on the last k_2 co-ordinates.

Then, there exists a constant $C = C(T) > 0$ such that for any (T, π_1, π_2) -family, \mathcal{A} , with $|\mathcal{A}| = n$, and an \mathcal{A} -closed set $S \subset \mathbb{R}^k$,

$$\sum_{i=0}^{k_2} b_i(\pi_3(S)) \leq C \cdot n^{(k_1+1)k_2}.$$

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Definition of cdcd

A cdcd of \mathbb{R}^k is a finite partition of \mathbb{R}^k into definable sets $(C_i)_{i \in I}$ (called the cells of the cdcd) satisfying the following properties. If $k = 1$ then a cdcd of \mathbb{R} is given by a finite set of points $a_1 < \dots < a_N$ and the cells of the cdcd are the singletons $\{a_i\}$ as well as the open intervals, $(\infty, a_1), (a_1, a_2), \dots, (a_N, \infty)$. If $k > 1$, then a cdcd of \mathbb{R}^k is given by a cdcd, $(C'_i)_{i \in I'}$, of \mathbb{R}^{k-1} and for each $i \in I'$, a collection of cells, C_i defined by,

$$C_i = \{\phi_i(C'_i \times D_j) \mid j \in J_i\},$$

Definition II

where

$$\phi_i : C'_i \times \mathbb{R} \rightarrow \mathbb{R}^k$$

is a definable homomorphism satisfying $\pi \circ \phi = \pi$, $(D_j)_{j \in J_i}$ is a cdcd of \mathbb{R} , and $\pi : \mathbb{R}^k \rightarrow \mathbb{R}^{k-1}$ is the projection map onto the first $k - 1$ coordinates. The cdcd of \mathbb{R}^k is then given by

$$\bigcup_{i \in I'} C_i.$$

Given a family of definable subsets $\mathcal{A} = \{S_1, \dots, S_n\}$ of \mathbb{R}^k , we say that a cdcd is adapted to \mathcal{A} , if each S_i is a union of cells of the given cdcd.

Quantitative cylindrical definable cell decomposition I

Theorem (Quantitative cylindrical definable cell decomposition)

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure over a real closed field \mathbb{R} , and let $T \subset \mathbb{R}^{k+\ell}$ be a closed definable set. Then, there exist constants $C_1, C_2 > 0$ depending only on T , and definable sets,

$$\{T_i\}_{i \in I}, \quad T_i \subset \mathbb{R}^k \times \mathbb{R}^{2(2^k-1)\cdot\ell},$$

depending only on T , with $|I| \leq C_1$, such that for any (T, π_1, π_2) -family, $\mathcal{A} = \{S_1, \dots, S_n\}$ with $S_i = T_{y_i}, y_i \in \mathbb{R}^\ell, 1 \leq i \leq n$, some sub-collection of the sets

Quantitative cylindrical definable cell decomposition II

Theorem (Quantitative cylindrical definable cell decomposition)

$$\pi_{k+2(2^k-1)\cdot\ell}^{\leq k} \left(\pi_{k+2(2^k-1)\cdot\ell}^{>k} \left(\mathbf{y}_{i_1}, \dots, \mathbf{y}_{i_{2(2^k-1)}} \right)^{-1} \cap T_i \right),$$
$$i \in I, 1 \leq i_1, \dots, i_{2(2^k-1)} \leq n,$$

form a cdcd of \mathbb{R}^k compatible with \mathcal{A} . Moreover, the cdcd has at most $C_2 \cdot n^{2(2^k-1)}$ cells.

Outline

- 1 Introduction
 - Some basic results
 - Combinatorial and Algebraic Complexity
- 2 Arrangements
- 3 O-minimal Structures and Admissible Sets
 - Examples of Admissible Sets
 - \mathcal{A} -sets
- 4 Results**
 - Bounds on Betti Numbers
 - Cylindrical Definable Cell Decomposition
 - Application: Generalization of a Theorem due to Alon et al.**
- 5 Idea of Proofs

Ramsey type theorem

Theorem

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure over a real closed field \mathbb{R} , and let F be a closed definable subset of $\mathbb{R}^\ell \times \mathbb{R}^\ell$. Then, there exists a constant $1 > \varepsilon = \varepsilon(F) > 0$, depending only on F , such that for any set of n points,

$$\mathcal{F} = \{\mathbf{y}_1, \dots, \mathbf{y}_n \in \mathbb{R}^\ell\}$$

there exists two subfamilies $\mathcal{F}_1, \mathcal{F}_2 \subset \mathcal{F}$, with $|\mathcal{F}_1|, |\mathcal{F}_2| \geq \varepsilon n$ and either,

- for all $\mathbf{y}_i \in \mathcal{F}_1$ and $\mathbf{y}_j \in \mathcal{F}_2$, $(\mathbf{y}_i, \mathbf{y}_j) \in F$, or
- for no $\mathbf{y}_i \in \mathcal{F}_1$ and $\mathbf{y}_j \in \mathcal{F}_2$, $(\mathbf{y}_i, \mathbf{y}_j) \in F$.

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Interesting corollary

Corollary

Let $\mathcal{S}(\mathbb{R})$ be an o-minimal structure over a real closed field \mathbb{R} , and let $T \subset \mathbb{R}^{k+\ell}$ be a closed definable set. Then, there exists a constant $1 > \varepsilon = \varepsilon(T) > 0$ depending only on T , such that for any (T, π_1, π_2) -family, $\mathcal{A} = \{S_1, \dots, S_n\}$, there exists two subfamilies $\mathcal{A}_1, \mathcal{A}_2 \subset \mathcal{A}$, with $|\mathcal{A}_1|, |\mathcal{A}_2| \geq \varepsilon n$, and either,

- for all $S_i \in \mathcal{A}_1$ and $S_j \in \mathcal{A}_2$, $S_i \cap S_j \neq \emptyset$, or
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Unions of definable families

Suppose that $T_1, \dots, T_m \subset \mathbb{R}^{k+l}$ are closed, definable sets, $\pi_1 : \mathbb{R}^{k+l} \rightarrow \mathbb{R}^k$ and $\pi_2 : \mathbb{R}^{k+l} \rightarrow \mathbb{R}^l$ the two projections.

Lemma

For any collection of (T_i, π_1, π_2) families \mathcal{A}_i , $1 \leq i \leq m$, the family $\bigcup_{1 \leq i \leq m} \mathcal{A}_i$ is a (T', π'_1, π'_2) family where,

$$T' = \bigcup_{i=1}^m T_i \times \{e_i\} \subset \mathbb{R}^{k+l+m},$$

with e_i the i -th standard basis vector in \mathbb{R}^m , and $\pi'_1 : \mathbb{R}^{k+l+m} \rightarrow \mathbb{R}^k$ and $\pi'_2 : \mathbb{R}^{k+l+m} \rightarrow \mathbb{R}^{l+m}$, the projections onto the first k and the last $l+m$ coordinates respectively.

Notations

Given closed definable sets $X \subset V \subset \mathbb{R}^k$, and $\varepsilon > 0$, we denote

$$\text{OT}(X, V, \varepsilon) = \{\mathbf{x} \in V \mid d_X(\mathbf{x}) < \varepsilon\},$$

$$\text{CT}(X, V, \varepsilon) = \{\mathbf{x} \in V \mid d_X(\mathbf{x}) \leq \varepsilon\},$$

$$\text{BT}(X, V, \varepsilon) = \{\mathbf{x} \in V \mid d_X(\mathbf{x}) = \varepsilon\},$$

and finally for $\varepsilon_1 > \varepsilon_2 > 0$ we define

$$\text{Ann}(X, V, \varepsilon_1, \varepsilon_2) = \{\mathbf{x} \in V \mid \varepsilon_2 < d_X(\mathbf{x}) < \varepsilon_1\},$$

$$\overline{\text{Ann}}(X, V, \varepsilon_1, \varepsilon_2) = \{\mathbf{x} \in V \mid \varepsilon_2 \leq d_X(\mathbf{x}) \leq \varepsilon_1\}.$$

Key Proposition

Proposition

Let $\mathcal{A} = \{S_1, \dots, S_n\}$ be a collection of closed definable subsets of \mathbb{R}^k and let $V \subset \mathbb{R}^k$ be a closed, and bounded definable set. Then, for all sufficiently small $1 \gg \varepsilon_1 \gg \varepsilon_2 > 0$ the following holds. For any connected component, C , of $\mathcal{A}(I) \cap V$, $I \subset [1 \dots n]$, there exists a connected component, D , of the definable set,

$$\bigcap_{1 \leq i \leq n} \text{Ann}(S_i, \varepsilon_1, \varepsilon_2)^c \cap V$$

such that D is definably homotopy equivalent to C .

Proof of Proposition

- For all sufficiently small $\varepsilon_1 > 0$ and for each connected component C of $\mathcal{A}(I) \cap V$, there exists a connected component D' of

$$\bigcap_{i \in I} S_i \cap \bigcap_{j \in [1 \dots n] \setminus I} \text{OT}(S_j, \varepsilon_1)^c \cap V,$$

homotopy equivalent to C .

- For $0 < \varepsilon_2 \ll \varepsilon_1$, and each connected component D' of $\bigcap_{i \in I} S_i \cap \bigcap_{j \in [1 \dots n] \setminus I} \text{OT}(S_j, \varepsilon_1)^c \cap V$, there exists a connected component D of $\bigcap_{i \in I} \text{CT}(S_i, \varepsilon_2) \cap \bigcap_{j \in [1 \dots n] \setminus I} \text{OT}(S_j, \varepsilon_1)^c \cap V$, homotopy equivalent to D' .

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Proof of Proposition (cont).

- Now notice that D is connected and contained in the set

$$\bigcap_{1 \leq i \leq n} \text{Ann}(\mathcal{S}_i, \varepsilon_1, \varepsilon_2)^c \cap V.$$

Let D'' be the connected component of

$$\bigcap_{1 \leq i \leq n} \text{Ann}(\mathcal{S}_i, \varepsilon_1, \varepsilon_2)^c \cap V$$

containing D .

- We claim that $D = D''$, which will prove the proposition.
- Suppose $D'' \setminus D \neq \emptyset$. Let $\mathbf{x} \in D'' \setminus D$ and \mathbf{y} any point in D . Since $\mathbf{x} \notin D$, either
 - there exists $i \in I$ such that $\mathbf{x} \in \text{OT}(\mathcal{S}_i, \varepsilon_1)^c$ or
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Proof of Proposition (cont).

- Let $\gamma : [0, 1] \rightarrow D''$ be a definable path with $\gamma(0) = \mathbf{x}$, $\gamma(1) = \mathbf{y}$. and let $d_i : D'' \rightarrow \mathbb{R}$ be the definable continuous function, $d_i(\mathbf{z}) = \text{dist}(\mathbf{z}, S_i)$.
- Then, in the first case, $d_i(\mathbf{x}) = d_i(\gamma(0)) \geq \varepsilon_1$ and $d_i(\mathbf{y}) = d_i(\gamma(1)) < \varepsilon_2$, implying that there exists $t \in (0, 1)$ with $\varepsilon_2 < d_i(\gamma(t)) < \varepsilon_1$ implying that $d_i(\gamma(t)) \notin \text{Ann}(S_i, \varepsilon_1, \varepsilon_2)^c$ and hence not in D'' (a contradiction).
- In the second case, $d_i(\mathbf{x}) = d_i(\gamma(0)) < \varepsilon_2$ and $d_i(\mathbf{y}) = d_i(\gamma(1)) \geq \varepsilon_1$, implying that there exists $t \in (0, 1)$ with $\varepsilon_2 < d_i(\gamma(t)) < \varepsilon_1$ again implying that $d_i(\gamma(t)) \notin \text{Ann}(S_i, \varepsilon_1, \varepsilon_2)^c$ and hence not in D'' (a contradiction).

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Proof of Theorem on Topological Complexity

- For $1 \leq i \leq n$, let $\mathbf{y}_i \in \mathbb{R}^\ell$ such that

$$S_i = T_{\mathbf{y}_i},$$

and let

$$A_i(\varepsilon_1, \varepsilon_2) = \text{Ann}(S_i, \varepsilon_1, \varepsilon_2)^c \cap V.$$

- Applying Mayer-Vietoris inequalities we have for $0 \leq i \leq k'$,

$$b_i\left(\bigcap_{j=1}^n A_j(\varepsilon_1, \varepsilon_2)\right) \leq b_{k'}(V) + \sum_{j=1}^{k'-i} \sum_{J \subset \{1, \dots, n\}, \#(J)=j} \left(b_{i+j-1}(A^J(\varepsilon_1, \varepsilon_2)) \right)$$

where $A^J(\varepsilon_1, \varepsilon_2) = \bigcup_{j \in J} A_j(\varepsilon_1, \varepsilon_2)$.

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Proof of Theorem on Topological Complexity (cont).

- Notice that each $\text{Ann}(S_i, \varepsilon_1, \varepsilon_2)^c$, $1 \leq i \leq n$, is a $(\text{Ann}(T, \varepsilon_1, \varepsilon_2)^c, \pi_1, \pi_2)$ -set and moreover,

$$\text{Ann}(S_i, \varepsilon_1, \varepsilon_2)^c = T_{y_i} \cap \text{Ann}(T, \varepsilon_1, \varepsilon_2)^c; \quad 1 \leq i \leq n.$$

- For $J \subset [1 \dots n]$, we denote

$$S^J(\varepsilon_1, \varepsilon_2) = \bigcup_{j \in J} \text{Ann}(S_j, \varepsilon_1, \varepsilon_2)^c.$$

There are only a finite number (depending on T) of topological types amongst $S^J(\varepsilon_1, \varepsilon_2)$. Restricting all the sets to V in the above argument, we obtain that there are only finitely many (depending on T and V) of topological types amongst the sets $A^J(\varepsilon_1, \varepsilon_2) = S^J(\varepsilon_1, \varepsilon_2) \cap V$.

Proof of Theorem on Topological Complexity (cont).

- Notice that each $\text{Ann}(\mathcal{S}_i, \varepsilon_1, \varepsilon_2)^c$, $1 \leq i \leq n$, is a $(\text{Ann}(T, \varepsilon_1, \varepsilon_2)^c, \pi_1, \pi_2)$ -set and moreover,

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There are only a finite number (depending on T) of topological types amongst $\mathcal{S}^J(\varepsilon_1, \varepsilon_2)$. Restricting all the sets to V in the above argument, we obtain that there are only finitely many (depending on T and V) of topological types amongst the sets $A^J(\varepsilon_1, \varepsilon_2) = \mathcal{S}^J(\varepsilon_1, \varepsilon_2) \cap V$.

Proof of Theorem on topological complexity(cont).

- Thus, there exists a constant $C(T, V)$ such that

$$C(T, V) = \max_{J \subset \{1, \dots, n\}} \left(b_{i+j-1}(A^J(\varepsilon_1, \varepsilon_2)) + b_{k'}(V) \right) + b_{k'}(V).$$

- It now follows from inequality ?? and Proposition 10 that,

$$\sum_{D \in \mathcal{C}(A, V)} b_i(D) \leq C \cdot n^{k'-i}.$$

Proof of Theorem on topological complexity(cont).

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Proof of Theorem for \mathcal{A} -sets

Key proposition:

Proposition

Let $\mathcal{A} = \{S_1, \dots, S_n\}$ be a collection of closed definable subsets of \mathbb{R}^k and let $V \subset \mathbb{R}^k$ be a closed, and bounded definable set and let S be an (\mathcal{A}, V) -closed set. Then, for all sufficiently small $1 \gg \varepsilon_1 \gg \varepsilon_2 \cdots \gg \varepsilon_n > 0$,

$$b(S) \leq \sum_{D \in \mathcal{C}(B, V)} b(D),$$

where

$$B = \bigcup_{i=1}^n \{S_i, \text{BT}(S_i, \varepsilon_i), \text{OT}(S_i, 2\varepsilon_i)^c\}.$$

Proof of Theorem on projections

Notice that for each $p, 0 \leq p \leq k_2$, and any \mathcal{A} -closed set $S \subset \mathbb{R}^{k_1+k_2}$, $W_{\pi_3}^p(S) \subset \mathbb{R}^{(p+1)k_1+k_2}$ is an \mathcal{A}^p -closed set where,

$$\mathcal{A}^p = \bigcup_{j=0}^p \mathcal{A}^{p,j},$$

$$\mathcal{A}^{p,j} = \bigcup_{i=1}^n \{S_i^{p,j}\},$$

where $S_i^{p,j} \subset \mathbb{R}^{(p+1)k_1+k_2}$ is defined by,

$$S_i^{p,j} = \{(\mathbf{x}_0, \dots, \mathbf{x}_p, \mathbf{y}) \mid \mathbf{x}_j \in \mathbb{R}^{k_1}, \mathbf{y} \in \mathbb{R}^{k_2}, (\mathbf{x}_j, \mathbf{y}) \in S_i\}.$$

and $W_f^i(X) = \{(\mathbf{x}_0, \dots, \mathbf{x}_i) \in X^{i+1} \mid f(\mathbf{x}_0) = \dots = f(\mathbf{x}_i)\}.$

Proof of Theorem on Projections (cont).

- Also, note that $\mathcal{A}^{p,j}$ is a $(T^{p,j}, \pi_1^p, \pi_2^p)$ family, where

$$T^{p,j} = \{(\mathbf{x}_0, \dots, \mathbf{x}_p, \mathbf{y}, \mathbf{z}) \mid \mathbf{x}_j \in \mathbb{R}^{k_1}, \mathbf{y} \in \mathbb{R}^{k_2}, \mathbf{z} \in \mathbb{R}^\ell, (\mathbf{x}_j, \mathbf{y}, \mathbf{z}) \in T, \text{ for some } j, 0 \leq j \leq p\}.$$

and $\pi_1^p : \mathbb{R}^{(p+1)k_1+k_2+\ell} \rightarrow \mathbb{R}^{(p+1)k_1+k_2}$, and $\pi_2^p : \mathbb{R}^{(p+1)k_1+k_2+\ell} \rightarrow \mathbb{R}^\ell$ are the appropriate projections.

- Since each $T^{p,j}$ is determined by T , we have using previous lemma that \mathcal{A}^p is a (T', π_1', π_2') -family for some definable T' determined by T .

Proof of Theorem on Projections (cont).

- Also, note that $\mathcal{A}^{p,j}$ is a $(T^{p,j}, \pi_1^p, \pi_2^p)$ family, where

$$T^{p,j} = \{(\mathbf{x}_0, \dots, \mathbf{x}_p, \mathbf{y}, \mathbf{z}) \mid \mathbf{x}_j \in \mathbb{R}^{k_1}, \mathbf{y} \in \mathbb{R}^{k_2}, \mathbf{z} \in \mathbb{R}^\ell, (\mathbf{x}_j, \mathbf{y}, \mathbf{z}) \in T, \text{ for some } j, 0 \leq j \leq p\}.$$

and $\pi_1^p : \mathbb{R}^{(p+1)k_1+k_2+\ell} \rightarrow \mathbb{R}^{(p+1)k_1+k_2}$, and $\pi_2^p : \mathbb{R}^{(p+1)k_1+k_2+\ell} \rightarrow \mathbb{R}^\ell$ are the appropriate projections.

- Since each $T^{p,j}$ is determined by T , we have using previous lemma that \mathcal{A}^p is a (T', π_1', π_2') -family for some definable T' determined by T .

Proof of Theorem on projections (cont).

Now $W_{\pi_3}^p(S) \subset \mathbb{R}^{(p+1)k_1+k_2}$ is a \mathcal{A}^p -closed set and $\#\mathcal{A}^p = (p+1)n$. Applying previous theorem we get, for each p and j , $0 \leq p, j < k_2$,

$$b_j(W_{\pi_3}^p(S)) \leq C_1(T) \cdot n^{(p+1)k_1+k_2}$$

The theorem now follows, since for each q , $0 \leq q < k_2$,

$$b_q(\pi_3(S)) \leq \sum_{i+j=q} b_j(W_{\pi_3}^i(S)) \leq C_2(T) \cdot n^{(q+1)k_1+k_2} \leq C(T) \cdot n^{(k_1+1)k_2}.$$

Proof of Ramsey type Theorem

- For each $i, 1 \leq i \leq n$, let

$$A_i = \pi_{2^\ell}^{\leq \ell}(\pi_{2^\ell}^{> \ell-1}(\mathbf{y}_i) \cap F),$$

and $\mathcal{G} = \{A_i \mid 1 \leq i \leq n\}$. Note that \mathcal{G} is a $(R, \pi_{2^\ell}^{\leq \ell}, \pi_{2^\ell}^{> \ell})$ -family.

- We now use the Clarkson-Shor random sampling technique (using Theorem on cdc instead of vertical decomposition). Applying Theorem on quantitative cdc to some sub-family $\mathcal{G}_0 \subset \mathcal{G}$ of cardinality r , we get a decomposition of \mathbb{R}^ℓ into at most $Cr^{2(2^\ell-1)} = r^{O(1)}$ definable cells, each of them defined by at most $2(2^\ell - 1) = O(1)$ of the \mathbf{y}_i 's.

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Proof of Ramesey type Theorem (cont).

- Let τ be a cell of the cdcd of \mathcal{G}_0 and let $G \in \mathcal{G}$. We say that G crosses τ if $G \cap \tau \neq \emptyset$ and $\tau \not\subseteq G$. The standard theory of random sampling now ensures that we can choose \mathcal{G}_0 such that each cell of the cdcd of \mathcal{G}_0 is “crossed” by no more than $\frac{c_1 n \log r}{r}$ elements of \mathcal{G} , where c_1 is a constant depending only on F .
- For each cell τ of the cdcd of \mathcal{G}_0 , let \mathcal{G}_τ denote the set of elements of \mathcal{G} which cross τ and let $\mathcal{F}_\tau = \mathcal{F} \cap \tau$.
- Since the total number of cells in the cdcd of \mathcal{G}_0 is bounded by $r^{O(1)}$, there must exist a cell τ such that,

$$|\mathcal{F}_\tau| \geq \frac{n}{r^{O(1)}}.$$

Now, every element of $\mathcal{G} \setminus \mathcal{G}_\tau$ either fully contains τ or is disjoint from it

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Proof of Ramsey type Theorem (cont).

- Setting $\alpha = \frac{1}{r^{O(1)}}$ and $\beta = \frac{1}{2}(1 - \frac{c_1 \log r}{r})$ we have that there exists a set $\mathcal{F}' = \mathcal{F}_\tau$ of cardinality at least αn , and a subset \mathcal{G}' of cardinality at least βn such that either each element of \mathcal{F}' is contained in every element of \mathcal{G}' , or no element of \mathcal{F}' is contained in any element of \mathcal{G}' .
- The proof is complete by taking $\mathcal{F}_1 = \mathcal{F}'$, and $\mathcal{F}_2 = \{\mathbf{y}_i \mid A_i \in \mathcal{G}'\}$ and $\varepsilon = \min(\alpha, \beta)$.

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